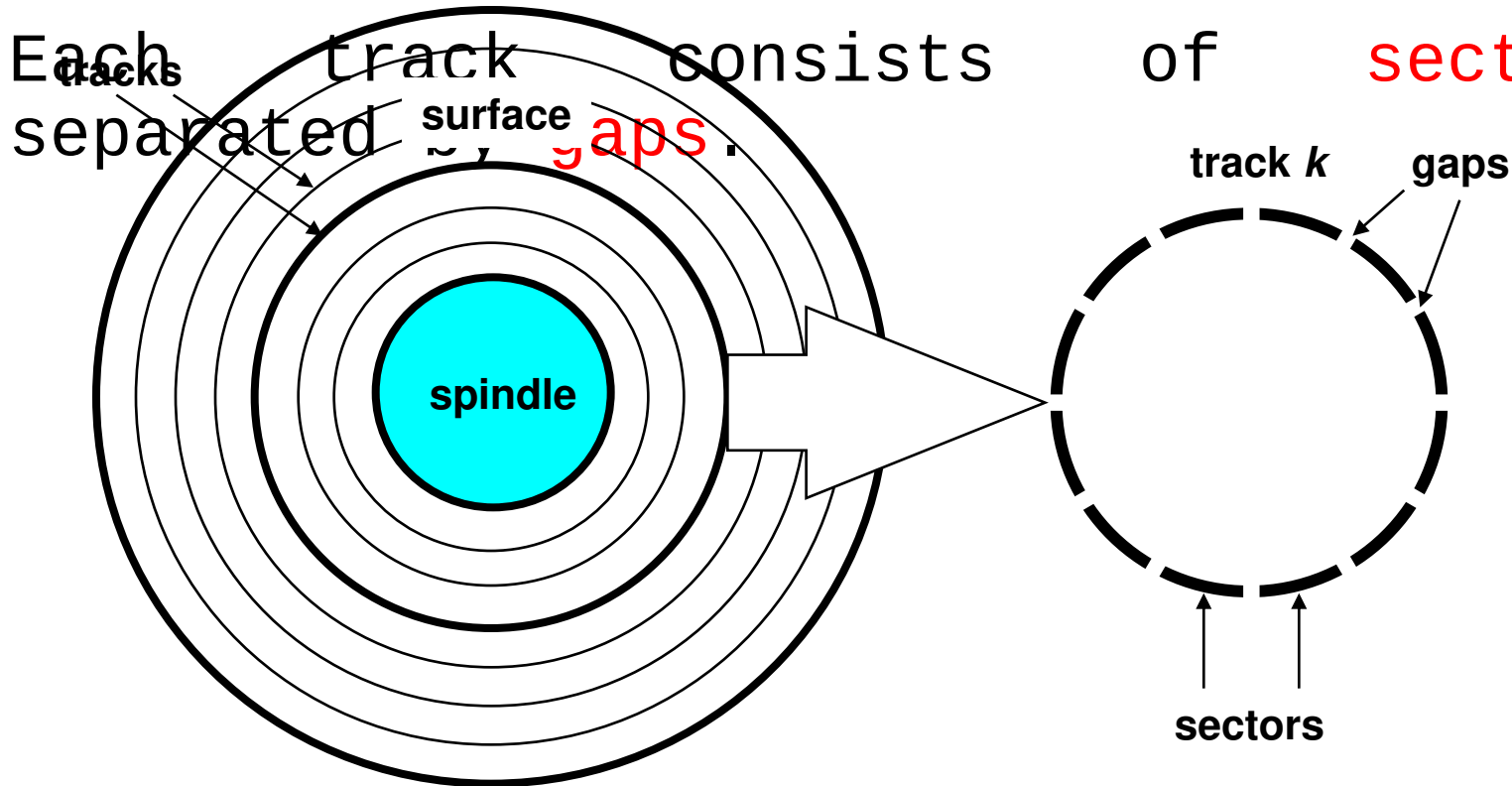


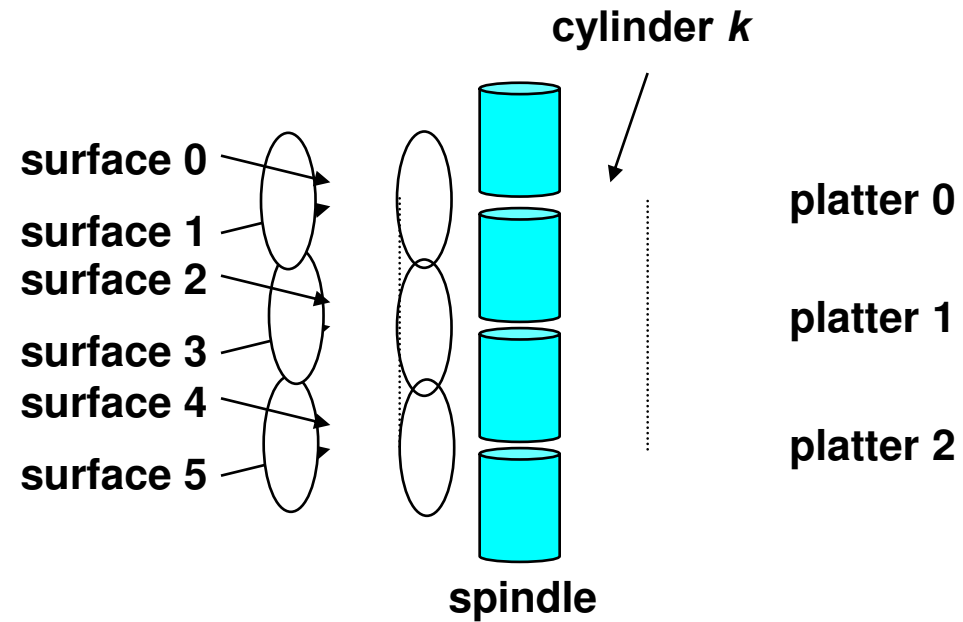
DISK GEOMETRY

- Disks consist of **platters**, each with two **surfaces**.
- Each surface consists of concentric rings called **tracks**.
- Each track consists of **sectors** separated by **gaps**.



DISK GEOMETRY (MULTIPLE-PLATTER VIEW)

- Aligned tracks form a cylinder.



DISK CAPACITY

- **Capacity**: maximum number of bits that can be stored.
 - Vendors express capacity in units of gigabytes (GB), where $1 \text{ GB} = 10^9$.
- Capacity is determined by these technology factors:
 - **Recording density** (bits/in): number of bits that can be squeezed into a 1 inch segment of a track.
 - **Track density** (tracks/in): number of tracks that can be squeezed into a 1 inch radial segment.
 - **Areal density** (bits/in²): product of recording and track density.

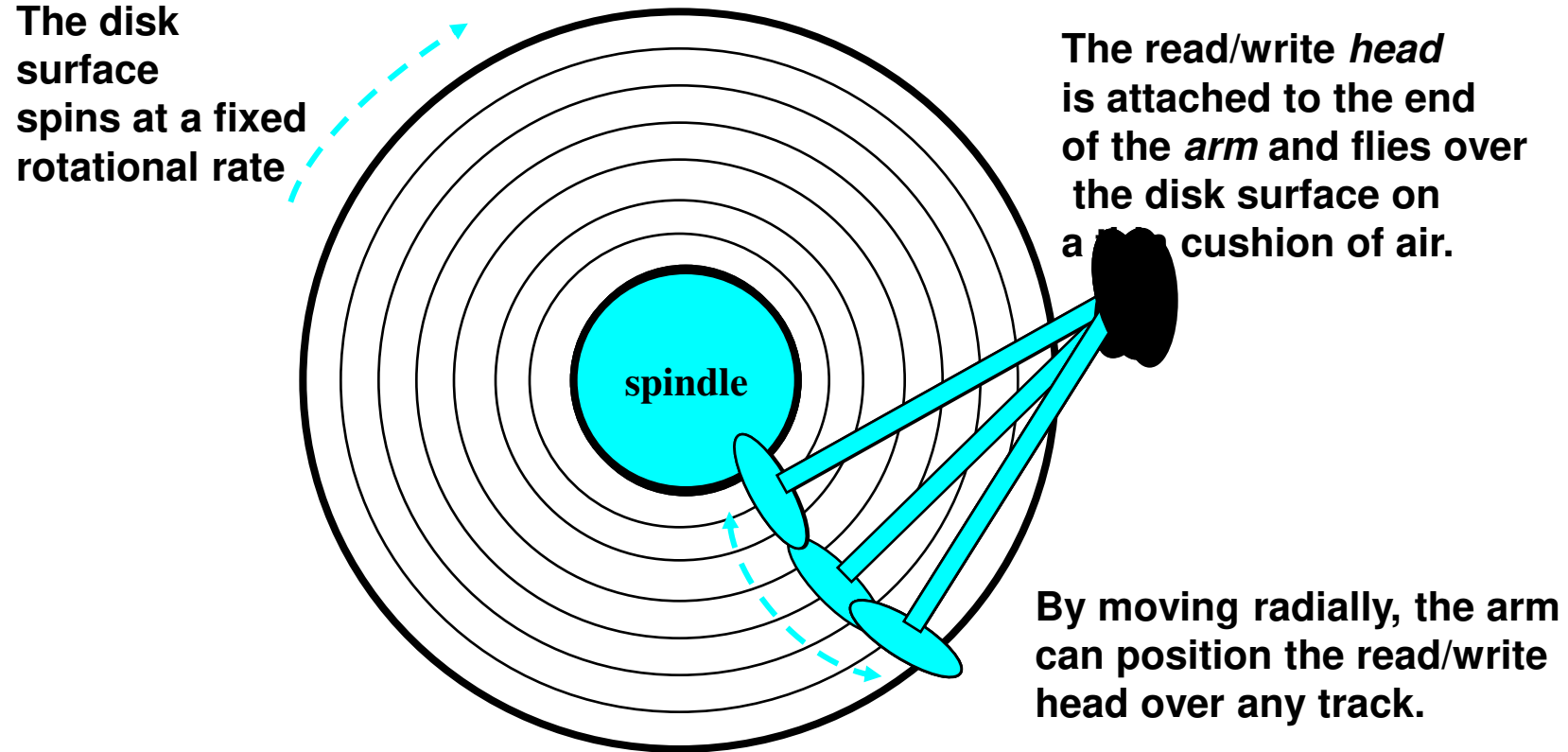
Cont...

- Modern disks partition tracks into disjoint subsets called **recording zones**
 - Each track in a zone has the same number of sectors, determined by the circumference of innermost track.
 - Each zone has a different number of sectors/track

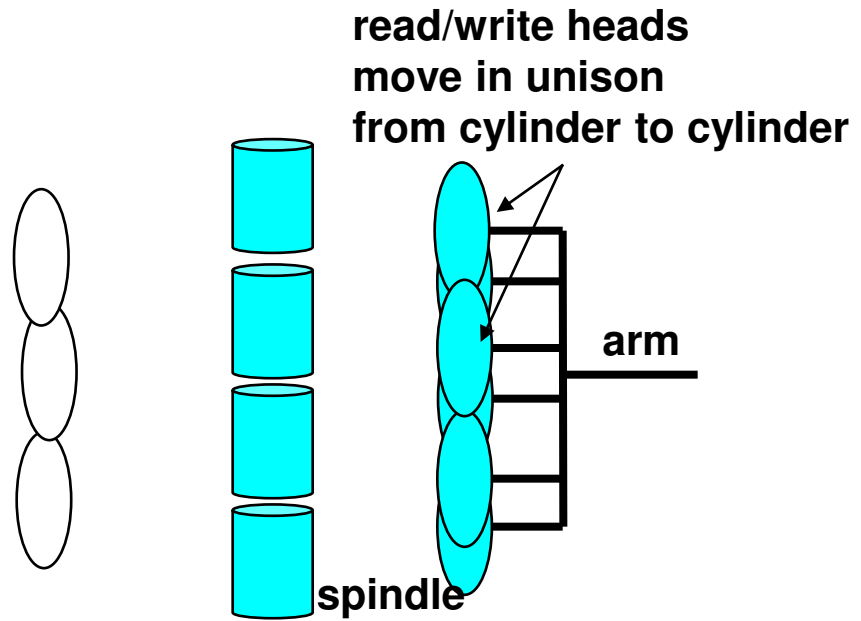
COMPUTING DISK CAPACITY

- Capacity = (# bytes/sector) x (avg. # sectors/track) x (# tracks/surface) x (# surfaces/platter) x (# platters/disk)
- Example:
 - 512 bytes/sector
 - 300 sectors/track (on average)
 - 20,000 tracks/surface
 - 2 surfaces/platter
 - 5 platters/disk
- Capacity = 512 x 300 x 20000 x 2 x 5
= 30,720,000,000 =
30.72 GB

DISK OPERATION (SINGLE-PLATTER VIEW)



DISK OPERATION (MULTI-PLATTER VIEW)



DISK ACCESS TIME

- Average time to access some target sector approximated by :
 - $T_{\text{access}} = T_{\text{avg seek}} + T_{\text{avg rotation}} + T_{\text{avg transfer}}$
- **Seek time** ($T_{\text{avg seek}}$)
 - Time to position heads over cylinder containing target sector.
 - Typical $T_{\text{avg seek}} = 9 \text{ ms}$
- **Rotational latency** ($T_{\text{avg rotation}}$)
 - Time waiting for first bit of target sector to pass under r/w head.
 - $T_{\text{avg rotation}} = 1/2 \times 1/\text{RPMs} \times 60 \text{ sec}/1 \text{ min}$

DISK ACCESS TIME

- **Transfer time** ($T_{\text{avg transfer}}$)
 - Time to read the bits in the target sector.
 - $T_{\text{avg transfer}} = 1/\text{RPM} \times 1/(\text{avg \# sectors/track}) \times 60 \text{ secs/1 min.}$

DISK ACCESS TIME EXAMPLE

- Given:

Rotational rate = 7,200 RPM

Average seek time = 9 ms.

Avg # sectors/track = 400.

- Derived:

$$T_{\text{avg rotation}} = \frac{1}{2} \times \left(\frac{60 \text{ secs}}{7200 \text{ RPM}} \right) \times 1000 \text{ ms/sec} = 4 \text{ ms.}$$

$$T_{\text{avg transfer}} = \frac{60}{7200} \text{ RPM} \times \frac{1}{400} \text{ secs/track} \times 1000 \text{ ms/sec} = 0.02 \text{ ms}$$

$$T_{\text{access}} = 9 \text{ ms} + 4 \text{ ms} + 0.02 \text{ ms}$$

DISK ACCESS TIME EXAMPLE

- Important points:
 - Access time dominated by seek time and rotational latency.
 - First bit in a sector is the most expensive, the rest are free.
 - SRAM access time is about 4 ns/double word, DRAM about 60 ns
 - Disk is about 40,000 times slower than SRAM,
 - 2,500 times slower than DRAM.

Types of External Memory

- Magnetic Disk
 - RAID
 - Removable
- Optical
 - CD-ROM
 - CD-Recordable (CD-R)
 - CD-R/W
 - DVD
- Magnetic Tape

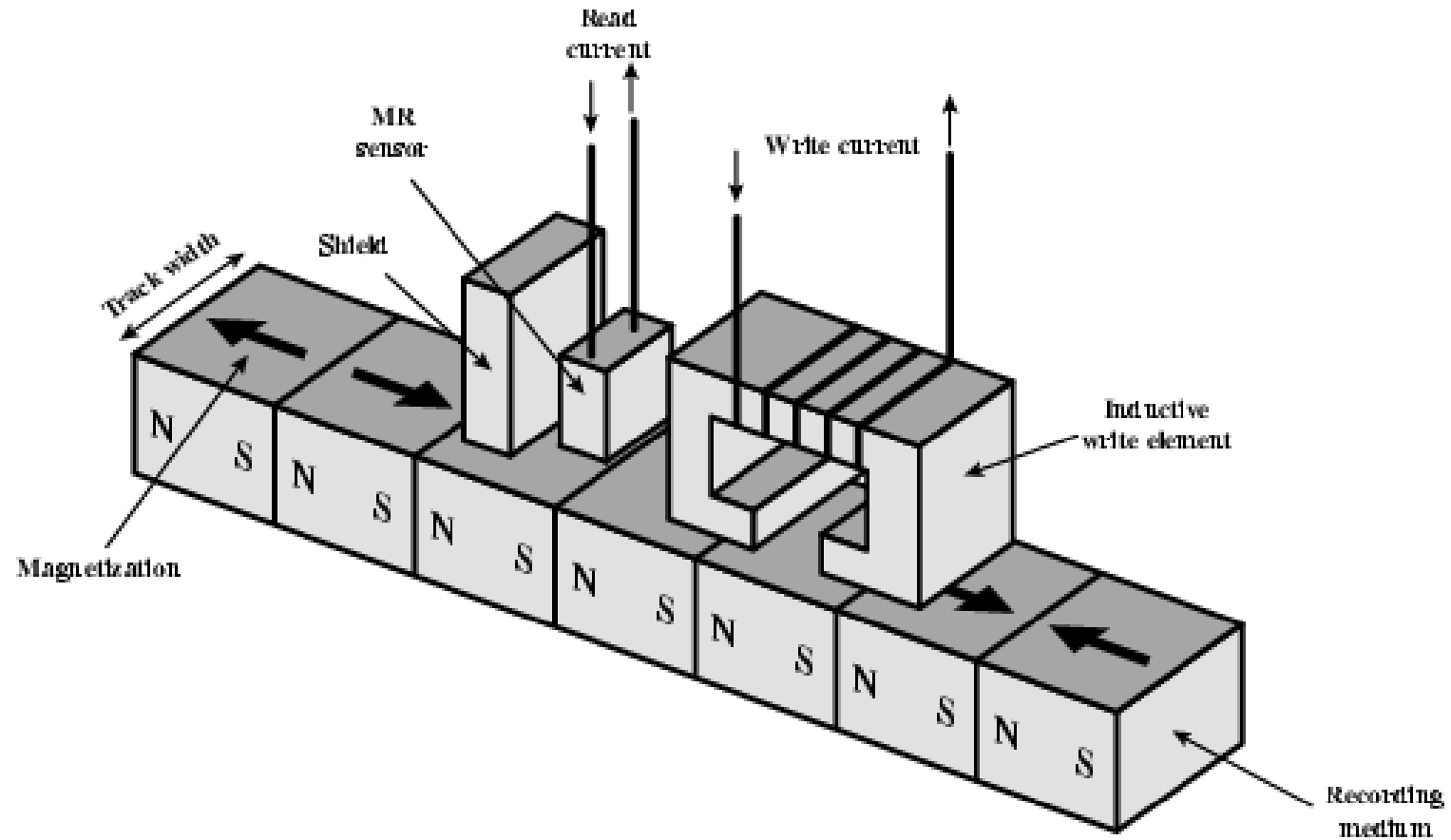
Magnetic Disk

- Disk substrate coated with magnetizable material (iron oxide...rust)
- Substrate used to be aluminium
- Now glass
 - Improved surface uniformity
 - Increases reliability
 - Reduction in surface defects
 - Reduced read/write errors
 - Lower flight heights (See later)
 - Better stiffness
 - Better shock/damage resistance

Read and Write Mechanisms

- Recording and retrieval via conductive coil called a head
- May be single read/write head or separate ones
- During read/write, head is stationary, platter rotates
- Write
 - Current through coil produces magnetic field
 - Pulses sent to head
 - Magnetic pattern recorded on surface below
- Read (traditional)
 - Magnetic field moving relative to coil produces current
 - Coil is the same for read and write
- Read (contemporary)
 - Separate read head, close to write head
 - Partially shielded magneto resistive (MR) sensor
 - Electrical resistance depends on direction of magnetic field
 - High frequency operation
 - Higher storage density and speed

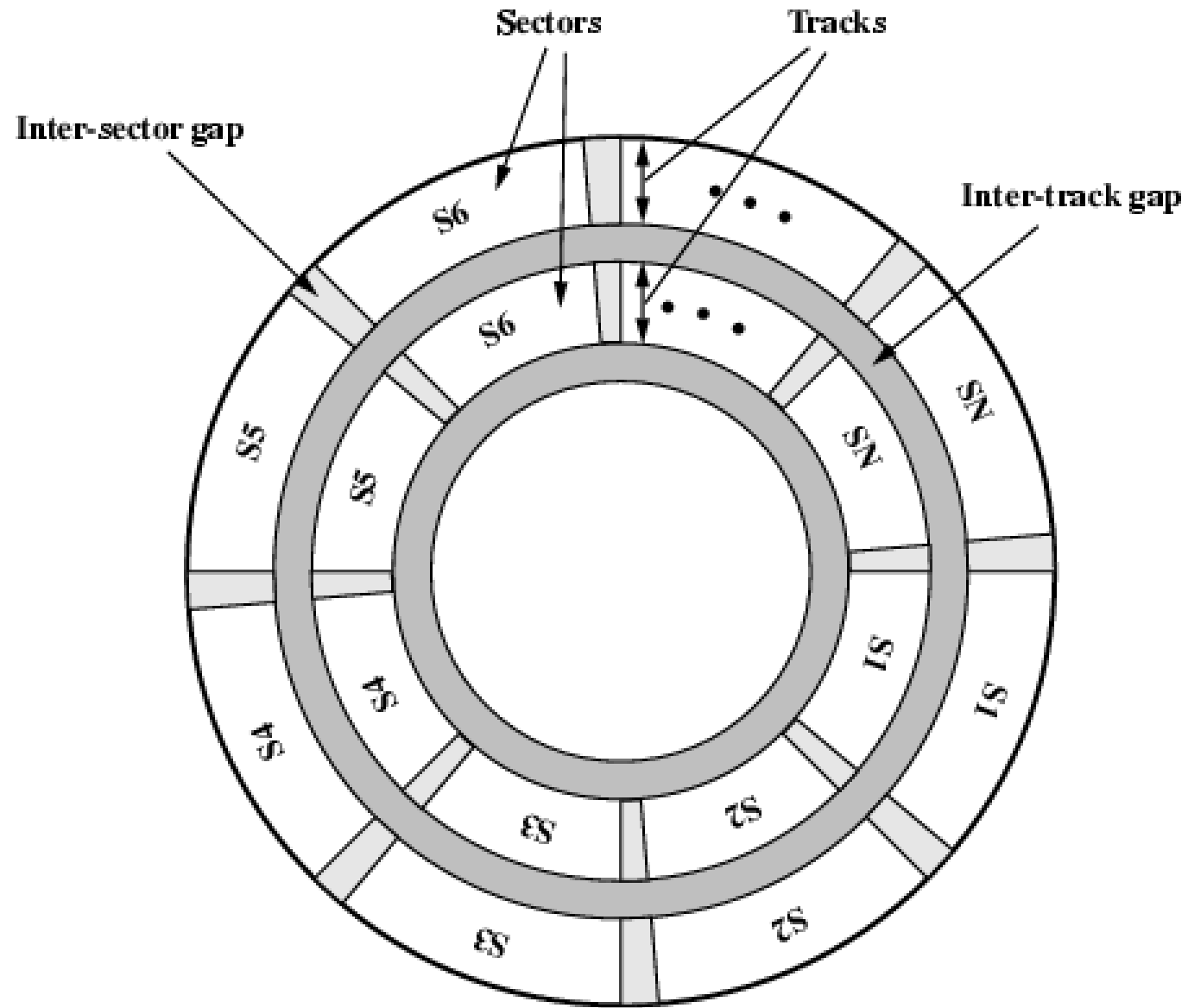
Inductive Write MR Read



Data Organization and Formatting

- Concentric rings or tracks
 - Gaps between tracks
 - Reduce gap to increase capacity
 - Same number of bits per track (variable packing density)
 - Constant angular velocity
- Tracks divided into sectors
- Minimum block size is one sector
- May have more than one sector per block

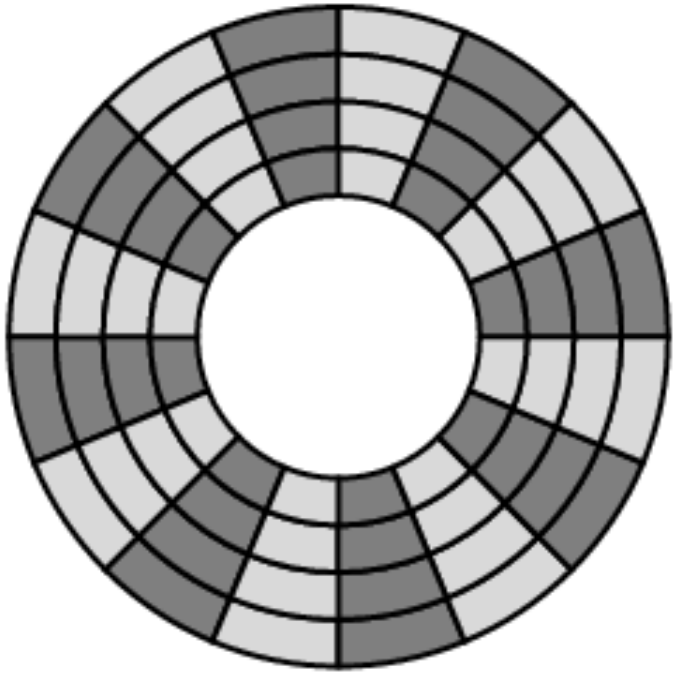
Disk Data Layout



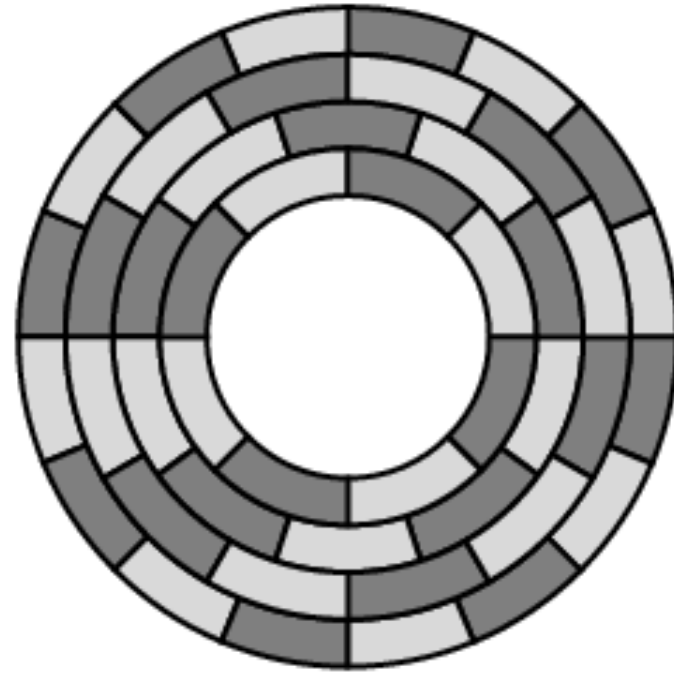
Disk Velocity

- Bit near centre of rotating disk passes fixed point slower than bit on outside of disk
- Increase spacing between bits in different tracks
- Rotate disk at constant angular velocity (CAV)
 - Gives pie shaped sectors and concentric tracks
 - Individual tracks and sectors addressable
 - Move head to given track and wait for given sector
 - Waste of space on outer tracks
 - Lower data density
- Can use zones to increase capacity
 - Each zone has fixed bits per track
 - More complex circuitry

Disk Layout Methods Diagram



(a) Constant angular velocity

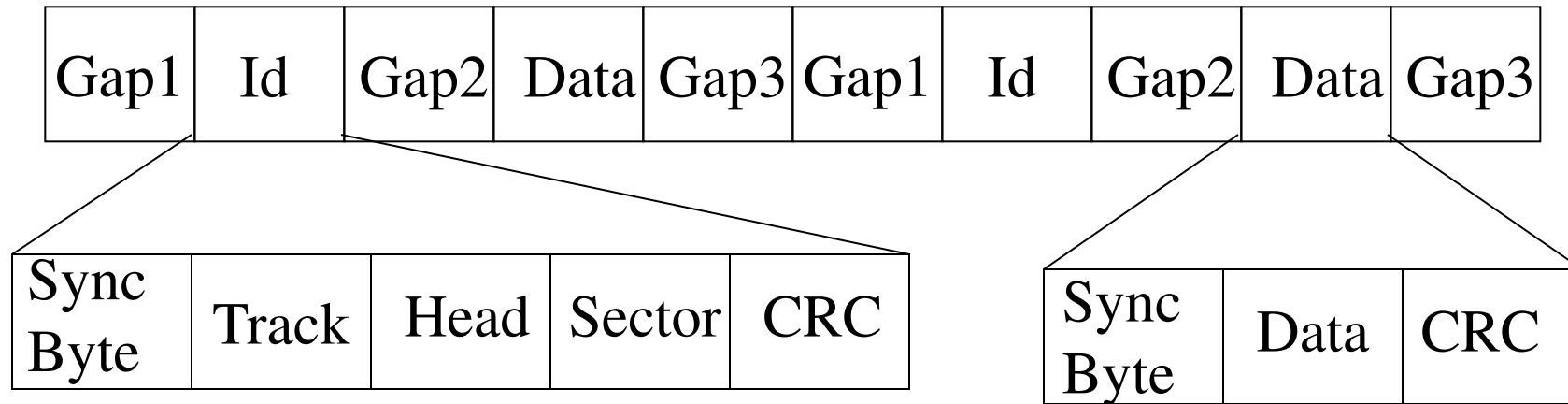


(b) Multiple zoned recording

Finding Sectors

- Must be able to identify start of track and sector
- Format disk
 - Additional information not available to user
 - Marks tracks and sectors

ST506 format (old!)



- Foreground reading
 - Find others

Characteristics

- Fixed (rare) or movable head
- Removable or fixed
- Single or double (usually) sided
- Single or multiple platter
- Head mechanism
 - Contact (Floppy)
 - Fixed gap
 - Flying (Winchester)

Fixed/Movable Head Disk

- Fixed head
 - One read write head per track
 - Heads mounted on fixed ridged arm
- Movable head
 - One read write head per side
 - Mounted on a movable arm

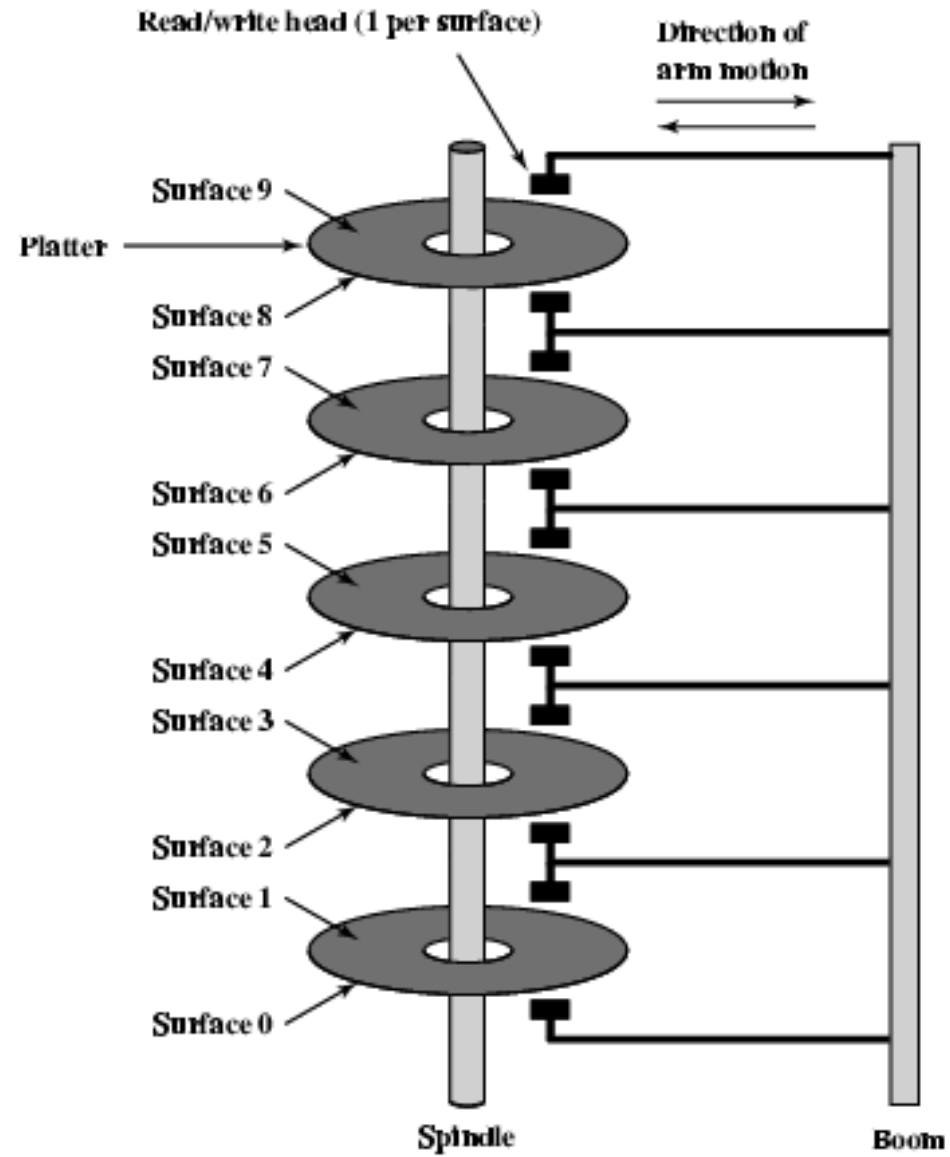
Removable or Not

- Removable disk
 - Can be removed from drive and replaced with another disk
 - Provides unlimited storage capacity
 - Easy data transfer between systems
- Nonremovable disk
 - Permanently mounted in the drive

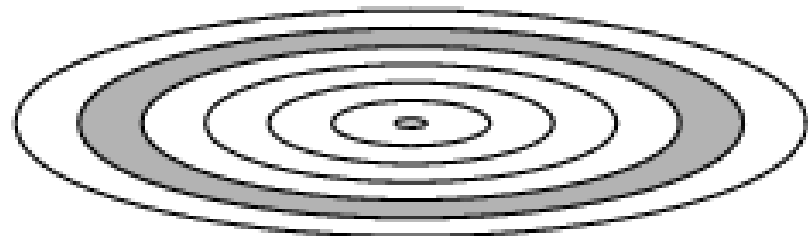
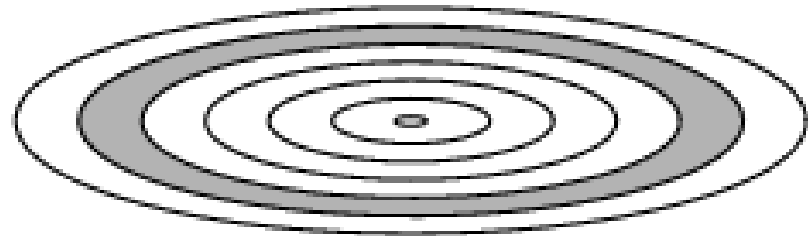
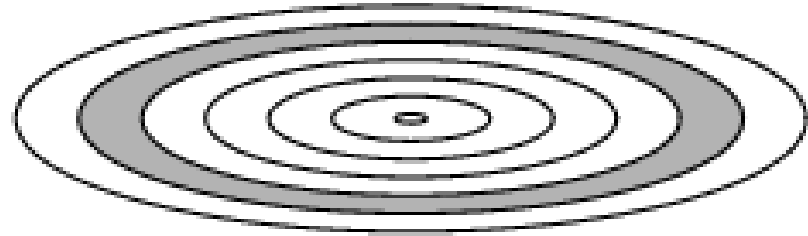
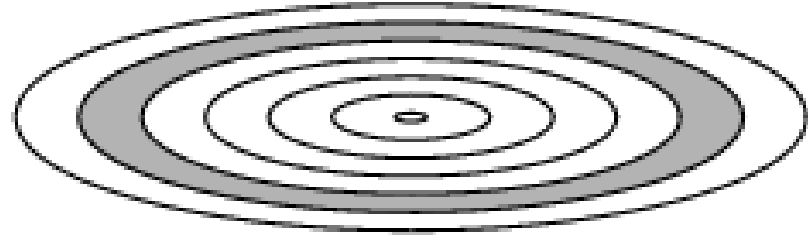
Multiple Platter

- One head per side
- Heads are joined and aligned
- Aligned tracks on each platter form cylinders
- Data is striped by cylinder
 - reduces head movement
 - Increases speed (transfer rate)

Multiple Platters



Cylinders



Floppy Disk

- 8", 5.25", 3.5"
- Small capacity
 - Up to 1.44Mbyte (2.88M never popular)
- Slow
- Universal
- Cheap
- Obsolete?

Winchester Hard Disk (1)

- Developed by IBM in Winchester (USA)
- Sealed unit
- One or more platters (disks)
- Heads fly on boundary layer of air as disk spins
- Very small head to disk gap
- Getting more robust

Winchester Hard Disk (2)

- Universal
- Cheap
- Fastest external storage
- Getting larger all the time
 - Multiple Gigabyte now usual

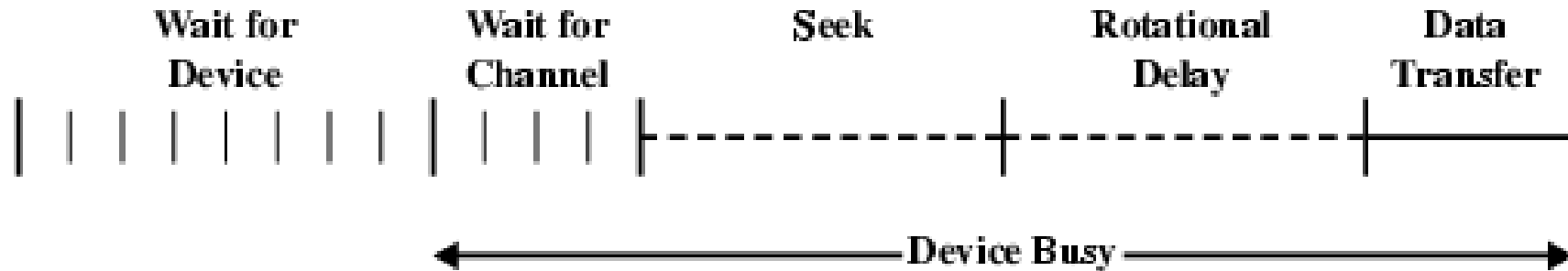
Removable Hard Disk

- ZIP
 - Cheap
 - Very common
 - Only 100M
- JAZ
 - Not cheap
 - 1G
- L-120 (a: drive)
 - Also reads 3.5" floppy
 - Becoming more popular?
- All obsoleted by CD-R and CD-R/W?

Speed

- Seek time
 - Moving head to correct track
- (Rotational) latency
 - Waiting for data to rotate under head
- Access time = Seek + Latency
- Transfer rate
- See timing comparison on p.173
 - Sequential organization = .064 seconds to read 1.28 MB
 - Random organization = 4.008 seconds to 1.28 MB
 - Why periodic defragmentation is important!

Timing of Disk I/O Transfer



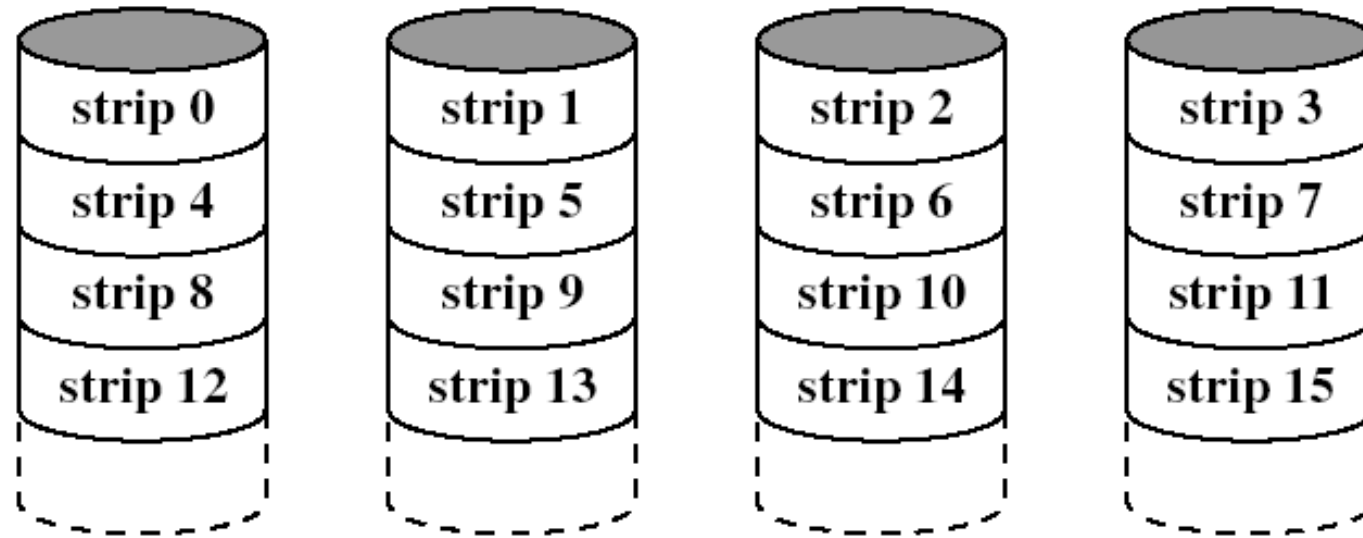
RAID

- Redundant Array of Independent Disks
- Redundant Array of Inexpensive Disks
- 6 levels in common use
- Not a hierarchy—but different design architectures with 3 common characteristics:
 - Set of physical disks viewed as single logical drive by O/S
 - Data distributed across physical drives of an array
 - Can use redundant capacity to store parity information, which guarantees data recoverability in case of a disk failure (except RAID 0)

RAID 0

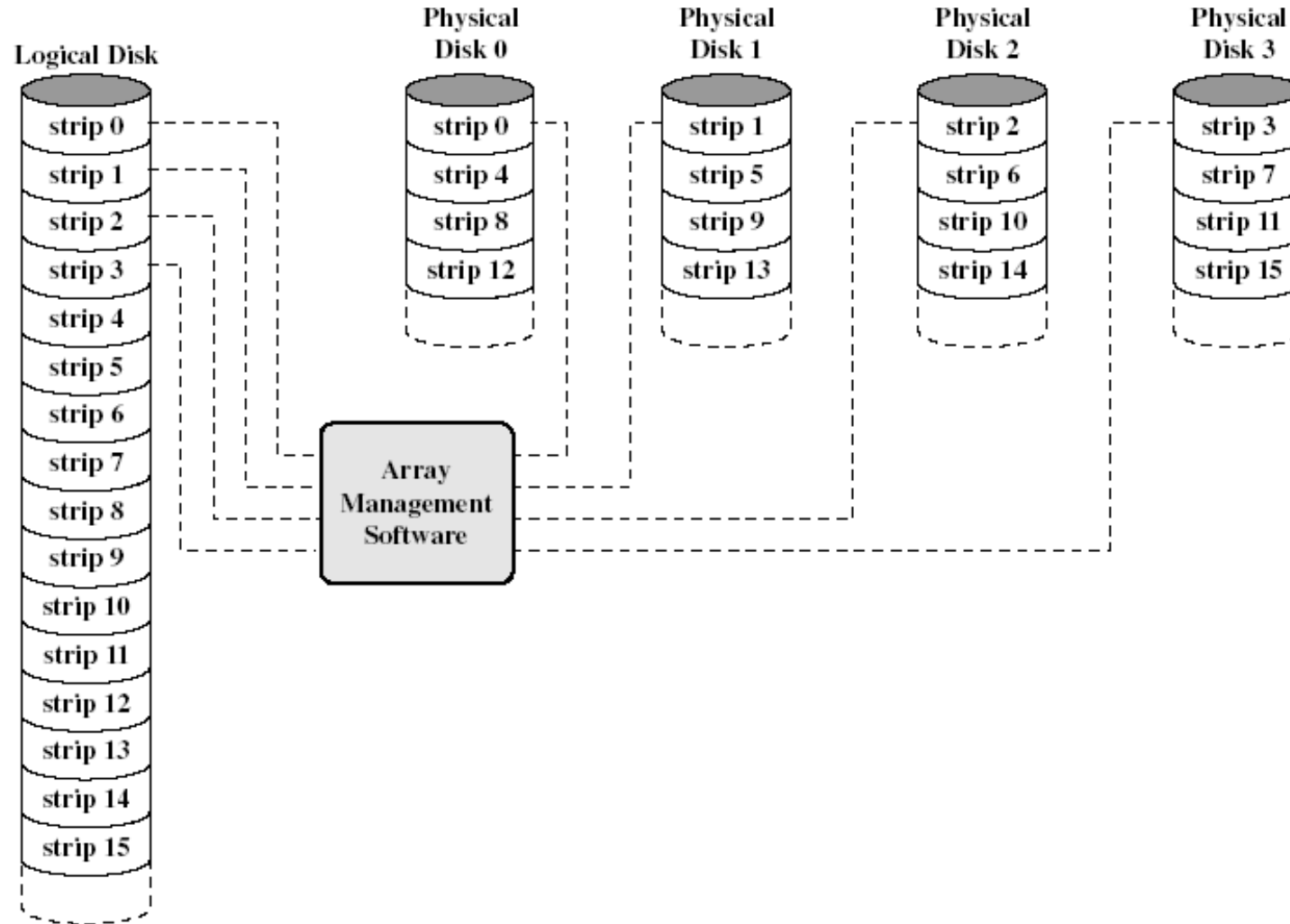
- No redundancy
- Data striped across all disks
- Round Robin striping
- Increase speed
 - Multiple data requests probably not on same disk
 - Disks seek in parallel
 - A set of data is likely to be striped across multiple disks
- **disk striping:** The procedure of combining a set of same-size disk partitions that reside on separate disks (from 2 to 32 disks) into a single volume, forming a virtual "stripe" across the disks that the operating system recognizes as a single drive. Disk striping enables multiple I/O operations in the same volume to proceed concurrently, thus offering enhanced performance. Microsoft Press« Computer and Internet Dictionary, 4th Edition, ©2000 Microsoft Corporation. All rights reserved.

RAID 0 Striping



(a) RAID 0 (non-redundant)

RAID 0 Logical Disk



RAID 1

- Mirrored Disks
- Data is striped across disks
- 2 copies of each stripe on separate disks
- Read from either to speed up reads
- Write to both at the same time in parallel
- Recovery is simple
 - Swap faulty disk & re-mirror
 - No down time
- Expensive—two complete sets of drives

RAID 2

- Disks are synchronized
- Very small stripes
 - Often single byte/word
- Error correction calculated across corresponding bits on disks
- *Multiple* parity disks store Hamming code error correction in corresponding positions
- Lots of redundancy
 - Expensive
 - Not commercially available

RAID 3

- Similar to RAID 2
- Only *one* redundant parity disk, no matter how large the array
- Simple parity bit for each set of corresponding bits
- Data on failed drive can be reconstructed from surviving data and parity info
- Very high transfer rates

RAID 4

- Each disk operates independently
- Good for high I/O request rate
- Large stripes
- Bit by bit parity calculated across stripes on each disk
- Parity stored on parity disk
- Not commercially available

RAID 5

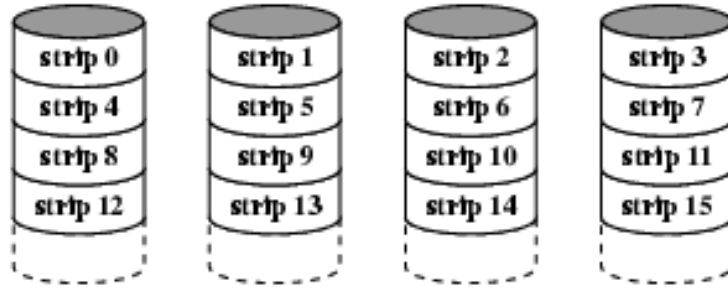
- Like RAID 4
- Parity striped across all disks
- Round robin allocation for parity stripe
- Avoids RAID 4 bottleneck at parity disk
- Commonly used in network servers

- N.B. DOES NOT MEAN 5 DISKS!!!!!!

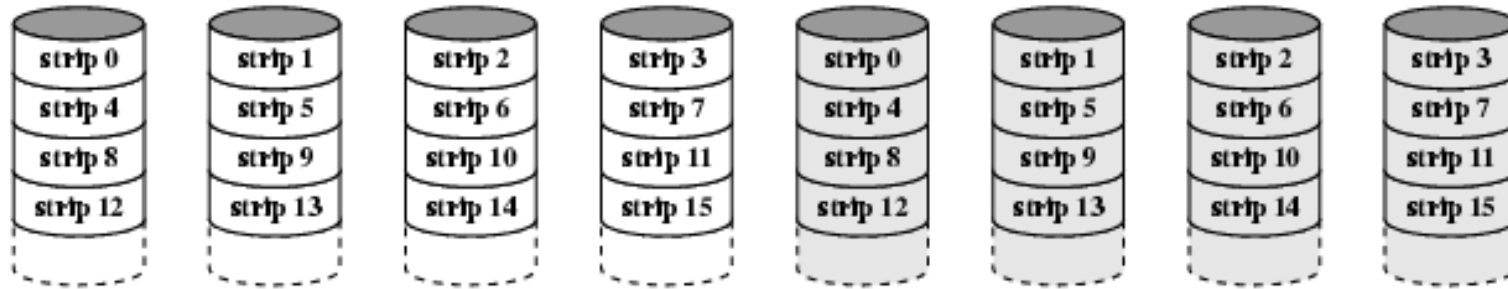
RAID 6

- Two parity calculations
- Stored in separate blocks on different disks
- User requirement of N disks needs $N+2$
- High data availability
 - Three disks need to fail for data loss
 - Significant write penalty

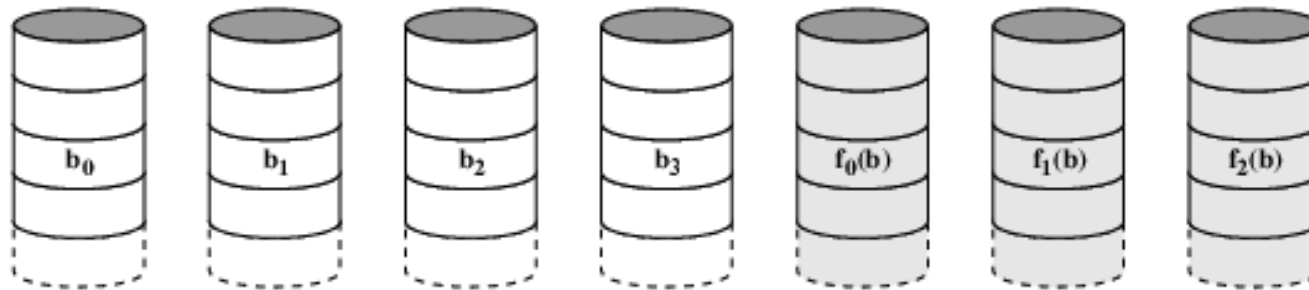
RAID 0, 1, 2



(a) RAID 0 (non-redundant)

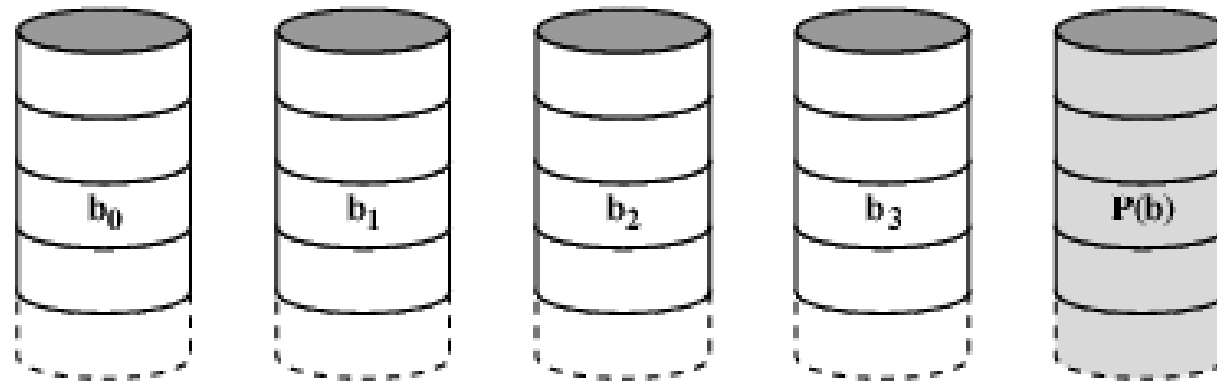


(b) RAID 1 (mirrored)

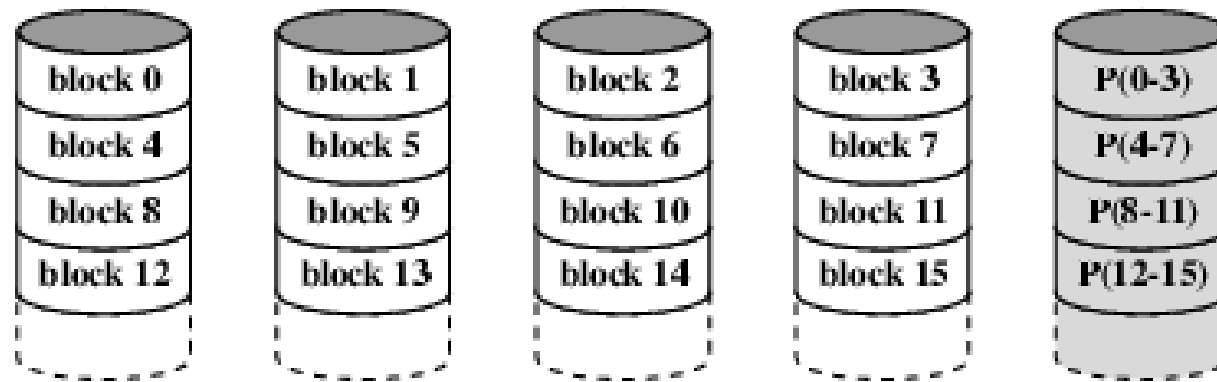


(c) RAID 2 (redundancy through Hamming code)

RAID 3 & 4

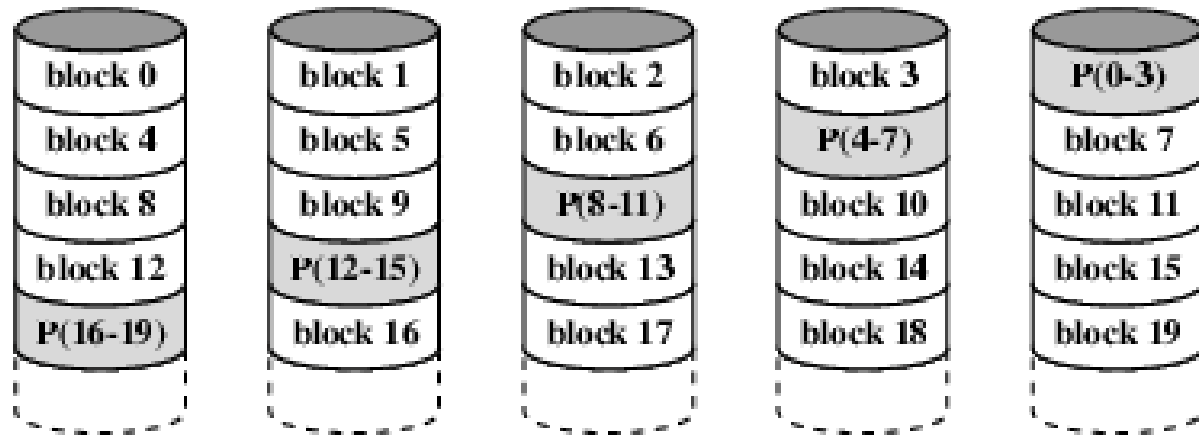


(d) RAID 3 (bit-interleaved parity)

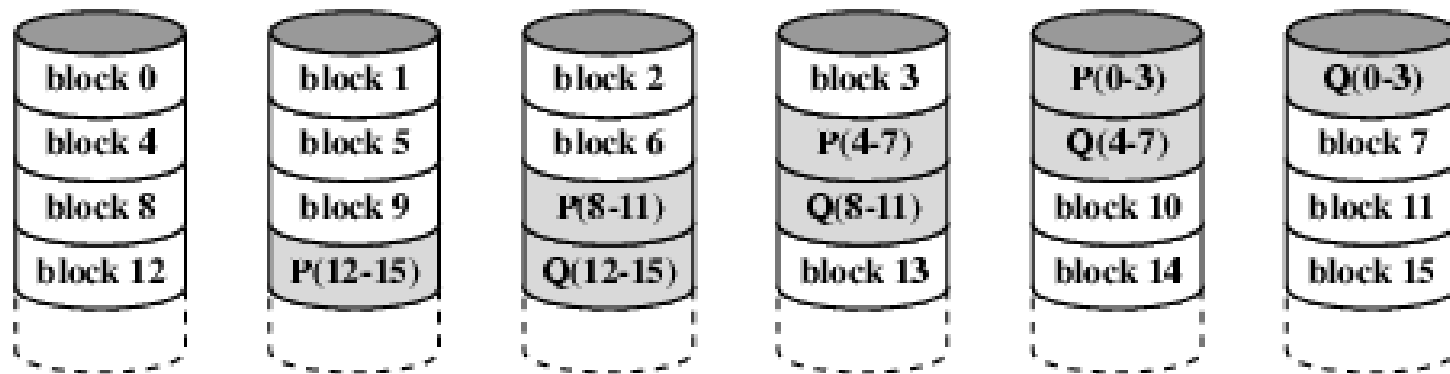


(e) RAID 4 (block-level parity)

RAID 5 & 6



(f) RAID 5 (block-level distributed parity)



(g) RAID 6 (dual redundancy)

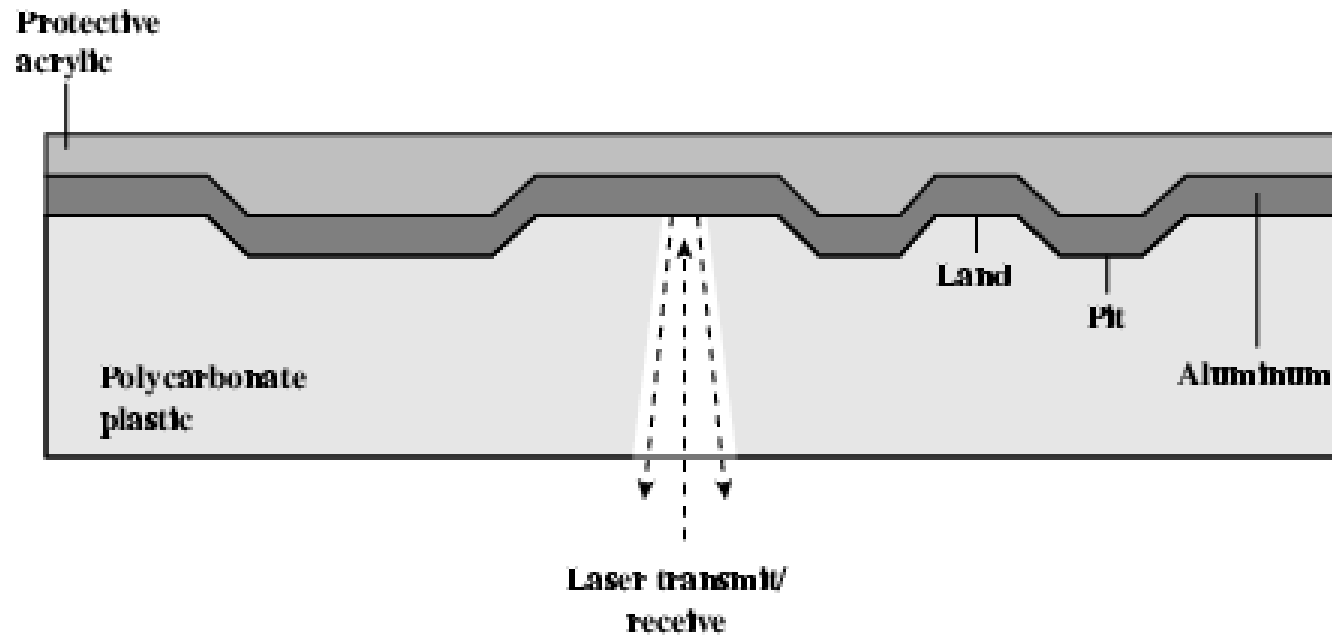
RAID Levels

| Category | Level | Description | I/O Request Rate (Read/Write) | Data Transfer Rate (Read/Write) | Typical Application |
|--------------------|-------|---|-------------------------------|---------------------------------|---|
| Striping | 0 | Non-redundant | Large strips: Excellent | Small strips: Excellent | Applications requiring high performance for non-critical data |
| Mirroring | 1 | Mirrored | Good/Fair | Fair/Fair | System drives; critical files |
| Parallel access | 2 | Redundant via Hamming code | Poor | Excellent | |
| | 3 | Bit-interleaved parity | Poor | Excellent | Large I/O request size applications, such as imaging, CAD |
| Independent access | 4 | Block-interleaved parity | Excellent/Fair | Fair/Poor | |
| | 5 | Block-interleaved distributed parity | Excellent/Fair | Fair/Poor | High request rate, read-intensive, data lookup |
| | 6 | Block-interleaved dual distributed parity | Excellent/Poor | Fair/Poor | Applications requiring extremely high availability |

Optical Storage CD-ROM

- Originally for audio
- 650Mbytes giving over 70 minutes audio
- Polycarbonate coated with highly reflective coat, usually aluminium
- Data stored as pits
- Read by reflecting laser
- Constant packing density
- Constant linear velocity

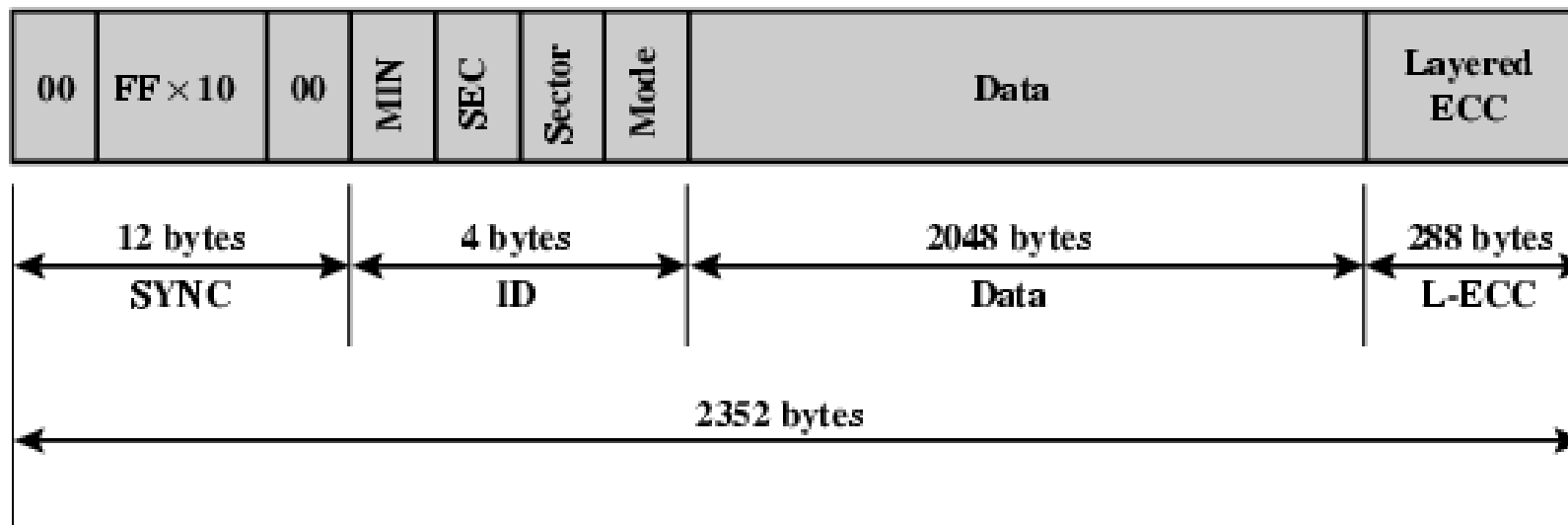
CD Operation



CD-ROM Drive Speeds

- Audio is single speed
 - Constant linear velocity
 - 1.2 ms^{-1}
 - Track (spiral) is 5.27km long
 - Gives 4391 seconds = 73.2 minutes
- Other speeds are quoted as multiples
- e.g. 24x
- Quoted figure is maximum drive can achieve

CD-ROM Format



- Mode 0=blank data field
- Mode 1=2048 byte data+error correction
- Mode 2=2336 byte data

Random Access on CD-ROM

- Difficult
- Move head to rough position
- Set correct speed
- Read address
- Adjust to required location
- (Yawn!)

CD-ROM for & against

- Large capacity (?)
 - Easy to mass produce
 - Removable
 - Robust
-
- Expensive for small runs
 - Slow
 - Read only

Other Optical Storage

- CD-Recordable (CD-R)
 - WORM
 - Now affordable
 - Compatible with CD-ROM drives
- CD-RW
 - Erasable
 - Getting cheaper
 - Mostly CD-ROM drive compatible
 - Phase change
 - Material has two different reflectivities in different phase states

DVD - what's in a name?

- Digital Video Disk
 - Used to indicate a player for movies
 - Only plays video disks
- Digital Versatile Disk
 - Used to indicate a computer drive
 - Will read computer disks and play video disks
- Dogs Veritable Dinner
- Officially - nothing!!!

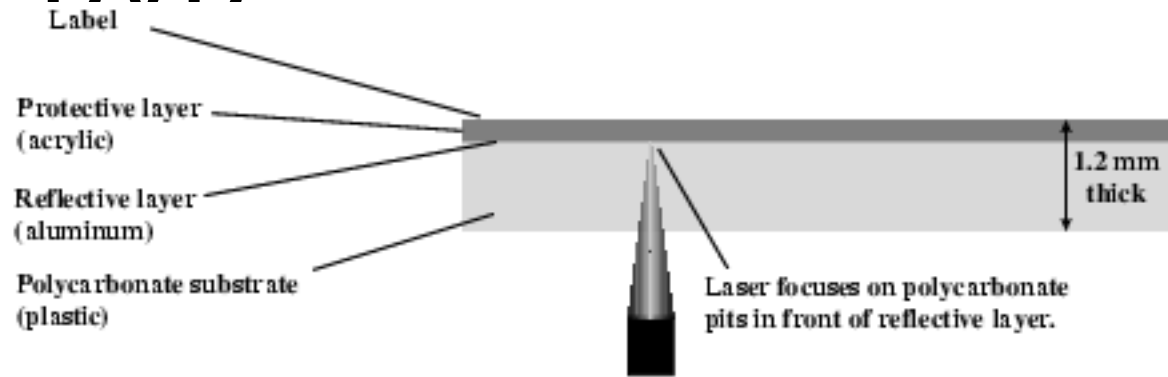
DVD - technology

- Multi-layer
- Very high capacity (4.7G per layer)
- Full length movie on single disk
 - Using MPEG compression
- Finally standardized (honest!)
- Movies carry regional coding
- Players only play correct region films
- Can be "fixed"

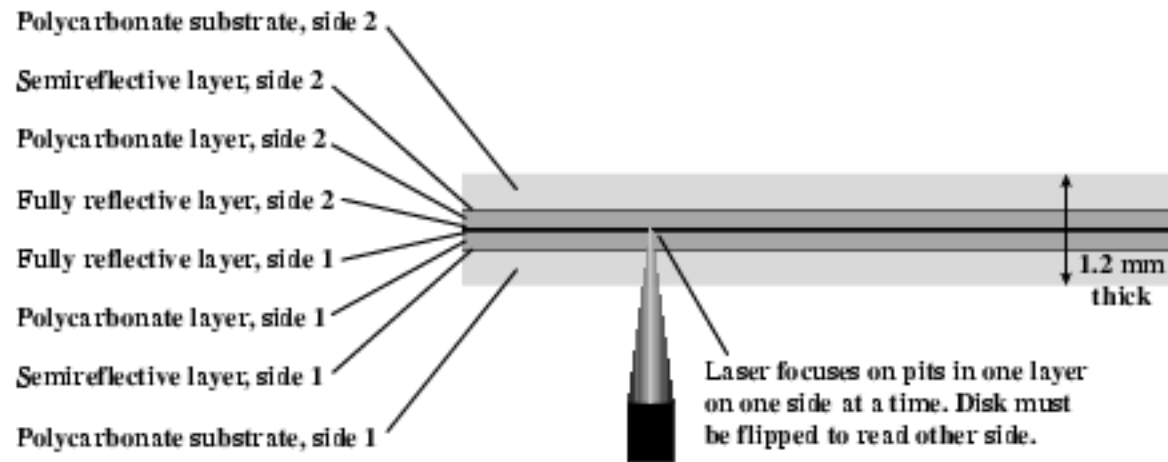
DVD – Writable

- Loads of trouble with standards
- First generation DVD drives may not read first generation DVD-W disks
- First generation DVD drives may not read CD-RW disks
- Wait for it to settle down before buying!

CD and DVD



(a) CD-ROM - Capacity 682 MB



(b) DVD-ROM, double-sided, dual-layer - Capacity 17 GB

Magnetic Tape

- Serial access
- Slow
- Very cheap
- Backup and archive

Digital Audio Tape (DAT)

- Uses rotating head (like video)
- High capacity on small tape
 - 4Gbyte uncompressed
 - 8Gbyte compressed
- Backup of PC/network servers