

Q1	Solve the Knapsack Problem using FIFOBB, assume knapsack capacity is $W = 8$ . Show how queue is used for node creation in the state space tree.	<table border="1"> <thead> <tr> <th>Item</th> <th>Profit</th> <th>Weight</th> </tr> </thead> <tbody> <tr> <td>1</td> <td>13</td> <td>4</td> </tr> <tr> <td>2</td> <td>15</td> <td>2</td> </tr> <tr> <td>3</td> <td>14</td> <td>4</td> </tr> <tr> <td>4</td> <td>16</td> <td>6</td> </tr> </tbody> </table>	Item	Profit	Weight	1	13	4	2	15	2	3	14	4	4	16	6
		Item	Profit	Weight													
		1	13	4													
		2	15	2													
		3	14	4													
4	16	6															

Answer

Ans:

Job 1, Job 2, Job 3 are selected with profit 12

**Node 1:**  $U = -(13 + 15) = -28$ , **Upper = -28**  
 $\hat{C} = -(28 + (8 - (4 + 2))14/4) = -35$ , Node 2, 3, 4, 5 created

**Node 2:**  $U = -(13 + 15) = -28$ ,  
 $\hat{C} = -(28 + (8 - (4 + 2))14/4) = -35$ , Node 6, 7, 8 created

**Node 3:**  $U = -(15 + 14) = -29$ , **Upper = -29**  
 $\hat{C} = -(29 + (8 - (2 + 4))16/2) = -41.33$ , Node 9, 10 created

**Node 4:**  $U = -(14) = -14$ , **Node killed  $U > \text{Upper}$**

**Node 5:**  $U = -(16) = -16$ , **Node killed  $U > \text{Upper}$**

**Node 6:**  $U = -(13 + 15) = -28$ , **Node killed  $U > \text{Upper}$**   
 $\hat{C} = -(28 + (8 - (4 + 2))14/4) = -35$ , Node 2, 3, 4, 5 created

**Node 7:**  $U = -(13 + 14) = -27$ , **Node killed  $U > \text{Upper}$**

**Node 8:**  $U = -(13) = -13$ , **Node killed  $U > \text{Upper}$**   
 $\hat{C} = -(29 + (8 - (2 + 4))16/2) = -41.33$ , Node 9, 10 created

**Node 9:**  $U = -(15 + 14) = -29$ ,  
 $\hat{C} = -(29 + (8 - (2 + 4))16/2) = -41.33$ , Node 11 created

**Node 10:**  $U = -(15 + 16) = -31$ , **Upper = -31**  
 $\hat{C} = -(31 + (8 - (2 + 6))) = -31$

**Node 11:**  $U = -(15 + 14) = -29$ , **Node killed  $U > \text{Upper}$**

Answer: Job2 and Job4 are selected with **Maximum profit = 31**

Q2	Find the existence of a pattern P in the given string S (assign digits A-C as 0-2), using Rabin Karp algorithm. For hash function use Mod 13. Find out how many spurious hits does the algorithm encounter in the <b>Text = ABCBBCABCBAABCCAACB</b> when looking for the pattern <b>Pattern = CCA</b> ?
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Answer	<p>Ans</p> $H(X) = \{P[2] \cdot 10^2 + P[1] \cdot 10^1 + P[0]\} \bmod 13.$ <p>Assign value <math>A=0, B=1, C=2</math></p> $H(P) = (200 + 20 + 0) \bmod 13 = 220 = 12.$
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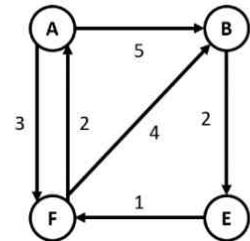
Text: ABCBBCABCBAABCCCAACB

ABC	12 MOD 13 =	12	spurious hits
BCB	121 MOD 13 =	4	
CBB	211 MOD 13 =	3	
BBC	112 MOD 13 =	8	
BCA	120 MOD 13 =	3	
CAB	201 MOD 13 =	6	
ABC	12 MOD 13 =	12	spurious hits
BCB	121 MOD 13 =	4	
CBA	210 MOD 13 =	2	
BAA	100 MOD 13 =	9	
AAA	0 MOD 13 =	0	
AAB	1 MOD 13 =	1	
ABC	12 MOD 13 =	12	spurious hits
BCC	122 MOD 13 =	5	
CCA	220 MOD 13 =	12	Hit
CAA	200 MOD 13 =	5	
AAC	2 MOD 13 =	2	
ACB	21 MOD 13 =	8	

Total number of spurious hits is 3 for the patten CCA.

Q3

Consider a logistics manager tasked with optimizing transportation routes for a delivery company that operates in a city with a complex network of roads. Your goal is to minimize the distance for packages to reach their destinations by finding the shortest paths between all pairs of locations. The transportation route is represented as weighted directed graph given below. Find the shortest paths between all pairs of locations, considering the varying distance, which helps company to delivery operation.



Answer

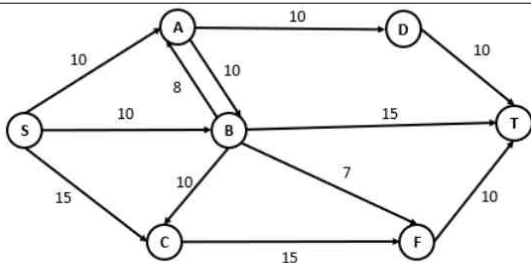
Base Condition (D0):

0 5 inf 3  
 inf 0 2 inf  
 inf inf 0 1  
 2 4 inf 0

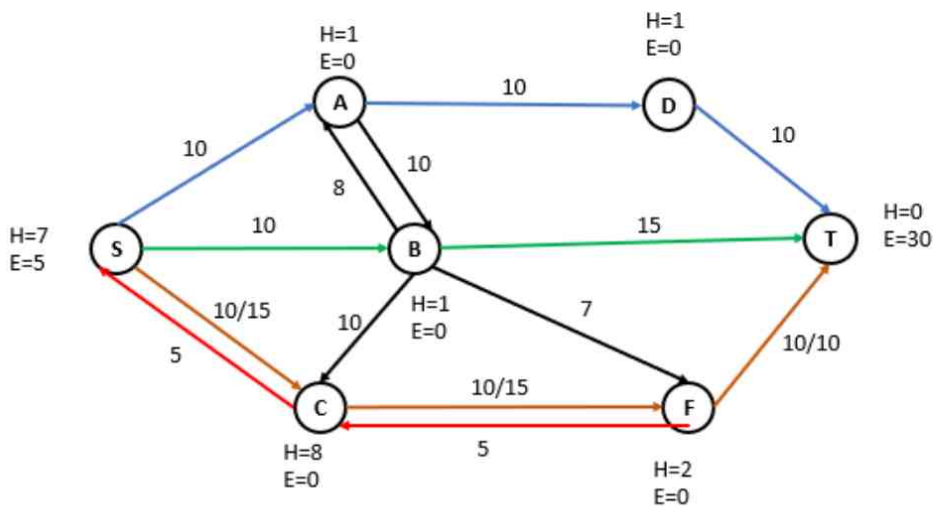
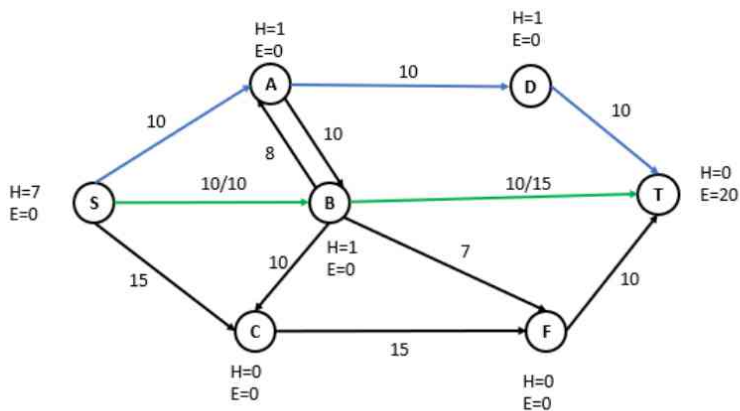
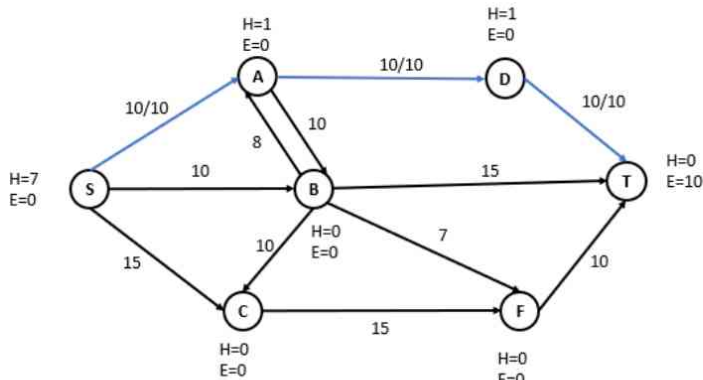
Iteration 1:	Shortest path from 0 to 0 with distance 0: 0
0 5 inf 3	Shortest path from 0 to 1 with distance 5: 0 -> 1
inf 0 2 inf	Shortest path from 0 to 3 with distance 3: 0 -> 3
inf inf 0 1	Shortest path from 1 to 1 with distance 0: 1
2 4 inf 0	Shortest path from 1 to 2 with distance 2: 1 -> 2
	Shortest path from 2 to 2 with distance 0: 2
	Shortest path from 2 to 3 with distance 1: 2 -> 3
	Shortest path from 3 to 0 with distance 2: 3 -> 0
	Shortest path from 3 to 1 with distance 4: 3 -> 1
	Shortest path from 3 to 3 with distance 0: 3
Iteration 2:	Shortest path from 0 to 0 with distance 0: 0
0 5 7 3	Shortest path from 0 to 1 with distance 5: 0 -> 1
inf 0 2 inf	Shortest path from 0 to 2 with distance 7: 0 -> 1 -> 2

<p>inf inf 0 1 2 4 6 0</p>	<p>Shortest path from 0 to 3 with distance 3: 0 -&gt; 3  Shortest path from 1 to 1 with distance 0: 1  Shortest path from 1 to 2 with distance 2: 1 -&gt; 2  Shortest path from 2 to 2 with distance 0: 2  Shortest path from 2 to 3 with distance 1: 2 -&gt; 3  Shortest path from 3 to 0 with distance 2: 3 -&gt; 0  Shortest path from 3 to 1 with distance 4: 3 -&gt; 1  Shortest path from 3 to 2 with distance 6: 3 -&gt; 1 -&gt; 2  Shortest path from 3 to 3 with distance 0: 3</p>
<p>Iteration 3: 0 5 7 3 inf 0 2 3 inf inf 0 1 2 4 6 0</p>	<p>Shortest path from 0 to 0 with distance 0: 0  Shortest path from 0 to 1 with distance 5: 0 -&gt; 1  Shortest path from 0 to 2 with distance 7: 0 -&gt; 1 -&gt; 2  Shortest path from 0 to 3 with distance 3: 0 -&gt; 3  Shortest path from 1 to 1 with distance 0: 1  Shortest path from 1 to 2 with distance 2: 1 -&gt; 2  Shortest path from 1 to 3 with distance 3: 1 -&gt; 2 -&gt; 3  Shortest path from 2 to 2 with distance 0: 2  Shortest path from 2 to 3 with distance 1: 2 -&gt; 3  Shortest path from 3 to 0 with distance 2: 3 -&gt; 0  Shortest path from 3 to 1 with distance 4: 3 -&gt; 1  Shortest path from 3 to 2 with distance 6: 3 -&gt; 1 -&gt; 2  Shortest path from 3 to 3 with distance 0: 3</p>
<p>Iteration 4: 0 5 7 3 5 0 2 3 3 5 0 1 2 4 6 0</p>	<p>Shortest path from 0 to 0 with distance 0: 0  Shortest path from 0 to 1 with distance 5: 0 -&gt; 1  Shortest path from 0 to 2 with distance 7: 0 -&gt; 1 -&gt; 2  Shortest path from 0 to 3 with distance 3: 0 -&gt; 3  Shortest path from 1 to 0 with distance 5: 1 -&gt; 2 -&gt; 3 -&gt; 0  Shortest path from 1 to 1 with distance 0: 1  Shortest path from 1 to 2 with distance 2: 1 -&gt; 2  Shortest path from 1 to 3 with distance 3: 1 -&gt; 2 -&gt; 3  Shortest path from 2 to 0 with distance 3: 2 -&gt; 3 -&gt; 0  Shortest path from 2 to 1 with distance 5: 2 -&gt; 3 -&gt; 1  Shortest path from 2 to 2 with distance 0: 2  Shortest path from 2 to 3 with distance 1: 2 -&gt; 3  Shortest path from 3 to 0 with distance 2: 3 -&gt; 0  Shortest path from 3 to 1 with distance 4: 3 -&gt; 1  Shortest path from 3 to 2 with distance 6: 3 -&gt; 1 -&gt; 2  Shortest path from 3 to 3 with distance 0: 3</p>

<p>Q4</p>	<p>In water distribution systems, we need to find the maximum amount of water that can be supplied from source S to destination T through a network pipes with capacity limitations. Given a directed graph <math>G=(V,E)</math> representing a water distribution system, where V is set of vertices and E is the set of edges, each edge (u,v) has a capacity <math>c(u,v)</math> representing the maximum water flow that can be supplied through the network pipe. Use Push Relabel algorithm to find the maximum water flow that can be supplied from node S to node T using given graph.</p>
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Answer: Ans: Max flow 30



- Q5 Find whether the following line segments intersect or not using cross product.
- $L1 : \{(1,23) \& (10,15)\}$  and  $L2 : \{(4,10) \& (6,20)\}$
  - $L3 : \{(4,5) \& (7,10)\}$  and  $L4 : \{(1,1) \& (5,5)\}$
  - $L5 : \{(1,1) \& (10,10)\}$  and  $L6 : \{(3,3) \& (5,5)\}$
  - $L7 : \{(1,1) \& (10,10)\}$  and  $L8 : \{(5,8) \& (3,3)\}$

Answer Ans:

a.  $P11 \times (\overline{P12 P21})$

$$= \begin{vmatrix} (4-1) & (10-1) \\ (10-23) & (15-23) \end{vmatrix} = \begin{vmatrix} 3 & 9 \\ -13 & -8 \end{vmatrix} = +ve$$

$$P11 \times (\overline{P12 P22})$$

$$= \begin{vmatrix} (6-1) & (10-1) \\ (20-23) & (15-23) \end{vmatrix} = \begin{vmatrix} 5 & 9 \\ -3 & -8 \end{vmatrix} = -ve$$

*L1 and L2 intersect*

b.  $P31 \times (\overline{P32 P41})$

$$= \begin{vmatrix} (1-4) & (7-4) \\ (1-5) & (10-5) \end{vmatrix} = \begin{vmatrix} -3 & 3 \\ -4 & 5 \end{vmatrix} = -ve$$

$$P31 \times (\overline{P32 P42})$$

$$= \begin{vmatrix} (5-4) & (7-4) \\ (5-5) & (10-5) \end{vmatrix} = \begin{vmatrix} 1 & 3 \\ 0 & 5 \end{vmatrix} = +ve$$

*L3 and L4 intersect*

c.  $P51 \times (\overline{P52 P61})$

$$= \begin{vmatrix} (3-1) & (10-1) \\ (3-1) & (10-1) \end{vmatrix} = \begin{vmatrix} 2 & 9 \\ 2 & 9 \end{vmatrix} = 0$$

$$P51 \times (\overline{P52 P62})$$

$$= \begin{vmatrix} (5-1) & (10-1) \\ (5-1) & (10-1) \end{vmatrix} = \begin{vmatrix} 4 & 9 \\ 4 & 9 \end{vmatrix} = 0$$

Bounding Box for L5

$$\widehat{P51} = \{\min(1, 10), \min(1, 10)\} = \{(1, 1)\}$$

$$\widehat{P52} = \{\max(1, 10), \max(1, 10)\} = \{(10, 10)\}$$

Bounding Box for L6

$$\widehat{P61} = \{\min(3, 5), \min(3, 5)\} = \{(3, 3)\}$$

$$\widehat{P62} = \{\max(3, 5), \max(3, 5)\} = \{(5, 5)\}$$

If  $\{P51, P52, \widehat{P51}, \widehat{P52}\}$  and  $\{P61, P62, \widehat{P61}, \widehat{P62}\}$

$(10 > 3)$  and  $(5 > 1)$  and  $(10 > 3)$  and  $(5 > 1)$  hence *L5 and L6 intersect*

d.  $P71 \times (\overline{P72 P81})$

$$= \begin{vmatrix} (5-1) & (10-1) \\ (8-1) & (10-1) \end{vmatrix} = \begin{vmatrix} 4 & 9 \\ 7 & 9 \end{vmatrix} = -ve$$

$$P71 \times (\overline{P72 P82})$$

$$= \begin{vmatrix} (3-1) & (10-1) \\ (3-1) & (10-1) \end{vmatrix} = \begin{vmatrix} 2 & 9 \\ 2 & 9 \end{vmatrix} = 0$$

*L1 and L2 intersect*

Line1 and Line2 points:

p1 (x y): 1 23

p2 (x y): 10 15

p3 (x y): 4 10

p4 (x y): 6 20

Position of P3 with respect to P1P2: right

Position of P4 with respect to P1P2: left

Position of P1 with respect to P3P4: left

Position of P2 with respect to P3P4: right

The line segments intersect.

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Line3 and Line4 points:

p1 (x y): 4 5

p2 (x y): 7 10

p3 (x y): 1 1

