

BCSE202L Data structures and Algorithms

Quicksort

Quicksort

- A

7	6	10	5	9	2	1	15	7
0	1	2	3	4	5	6	7	8

```
1  def partition(a, l, h):
2      pivot = a[l]
3      i = l
4      j = h
5      while i < j:
6          while a[i] <= pivot:
7              i = i + 1
8          while a[j] > pivot:
9              j = j - 1
10         if i < j:
11             swap(a[i], a[j])
12     swap(a[l], a[j])
13     return j
```

```
1 def quicksort(a, l, h):  
2     if l < h:  
3         pivot_index = prtition(a, l, h)  
4         quicksort(a, l, pivot_index-1)  
5         quicksort(a, pivot_index, u)
```

Example

We are given array of n integers to sort:

40	20	10	80	60	50	7	30	100
----	----	----	----	----	----	---	----	-----

Pick Pivot Element

There are a number of ways to pick the pivot element. In this example, we will use the first element in the array:

40	20	10	80	60	50	7	30	100
----	----	----	----	----	----	---	----	-----

Partitioning Array

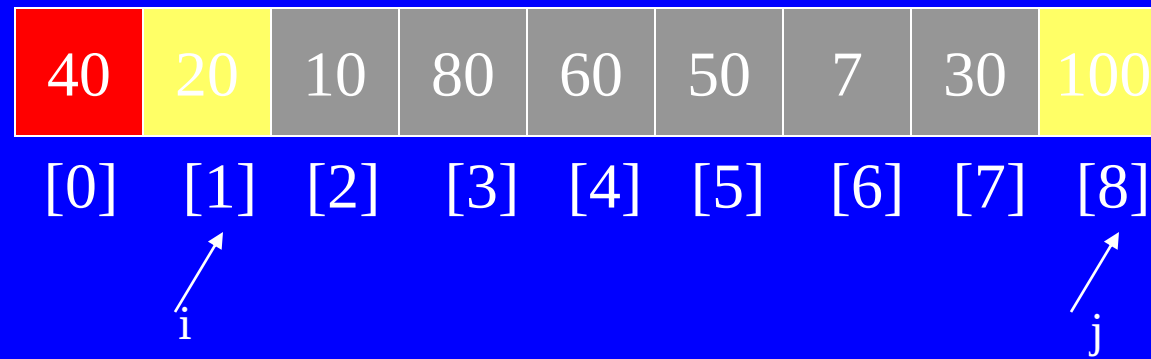
Given a pivot, partition the elements of the array such that the resulting array consists of:

1. One sub-array that contains elements $\geq$ pivot
2. Another sub-array that contains elements $<$ pivot

The sub-arrays are stored in the original data array.

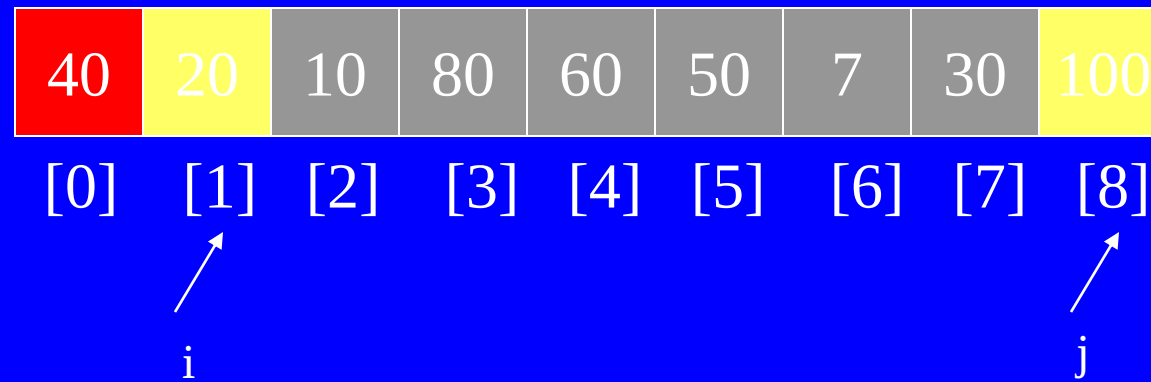
Partitioning loops through, swapping elements below/above pivot.

pivot = 40



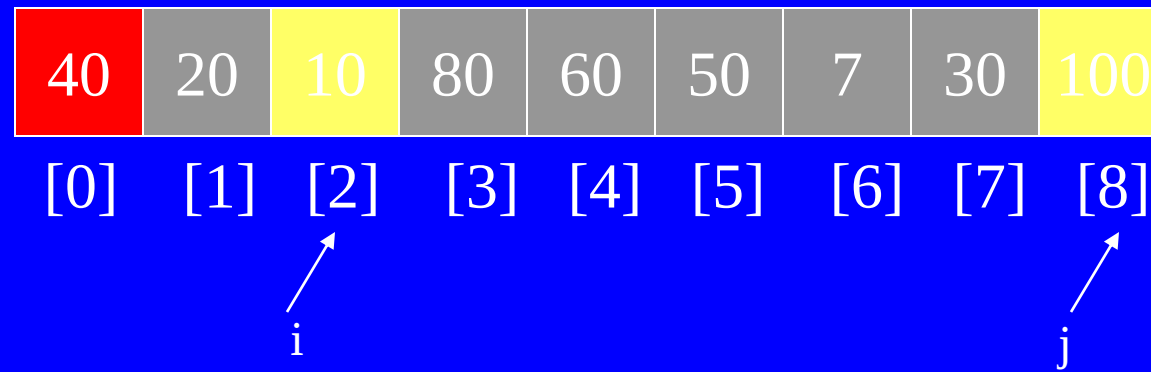
1. While $A[i] \leq \text{pivot}$
 $i++$

pivot = 40



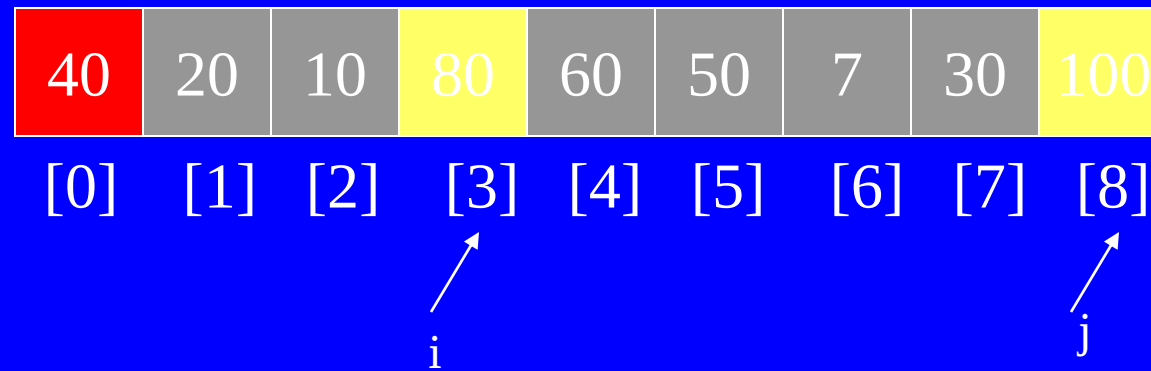
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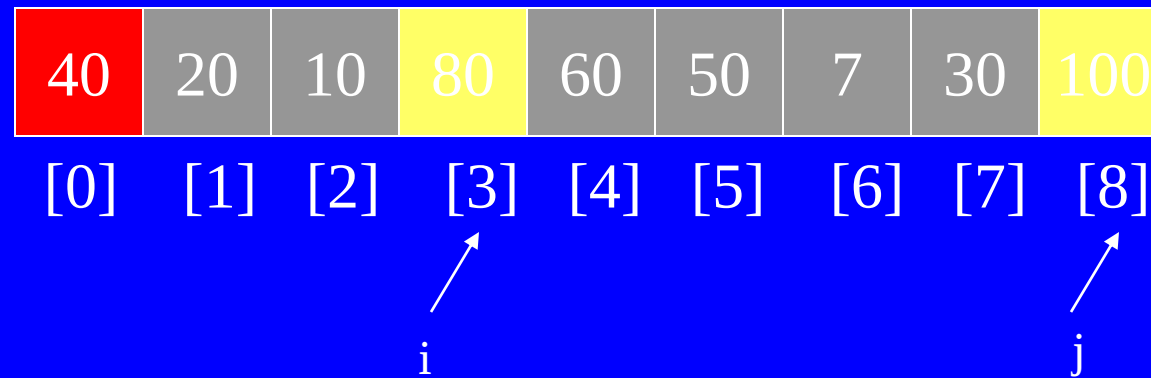
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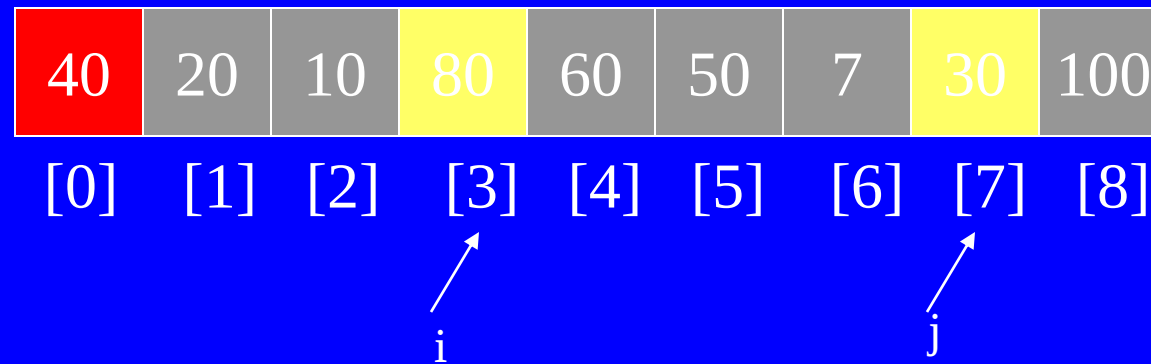
1. While $A[i] \leq \text{pivot}$
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2. While $A[j] > \text{pivot}$
 $j--$

pivot = 40



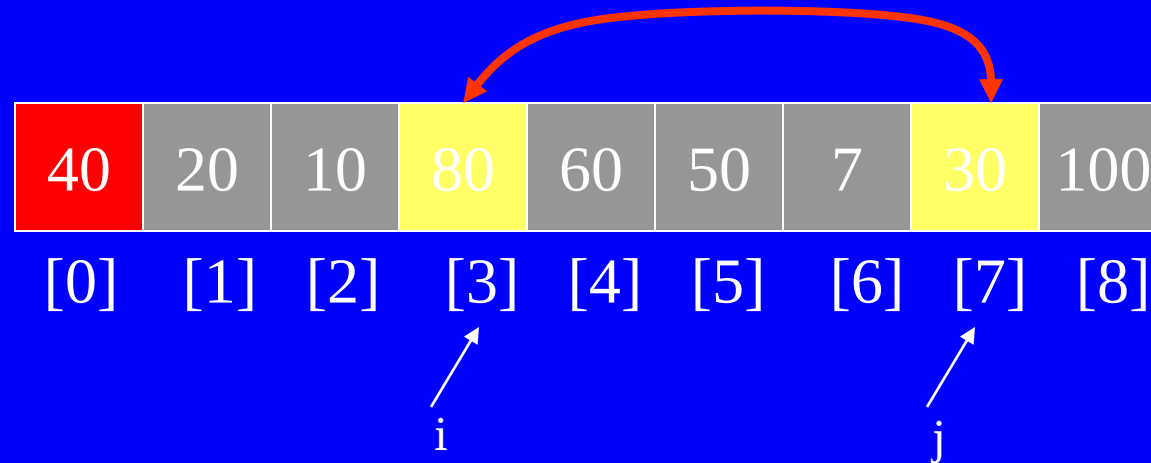
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$\text{pivot} = 40$



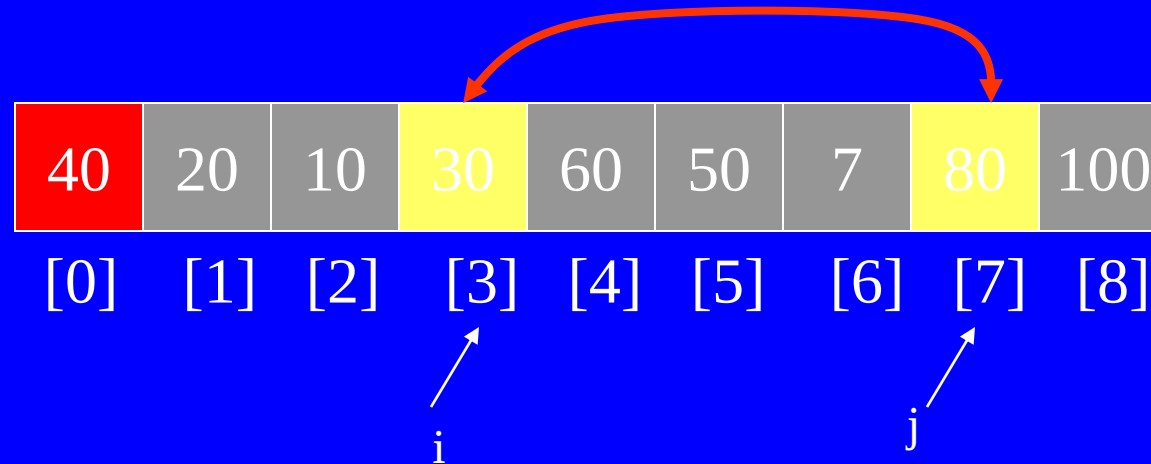
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			i				j	

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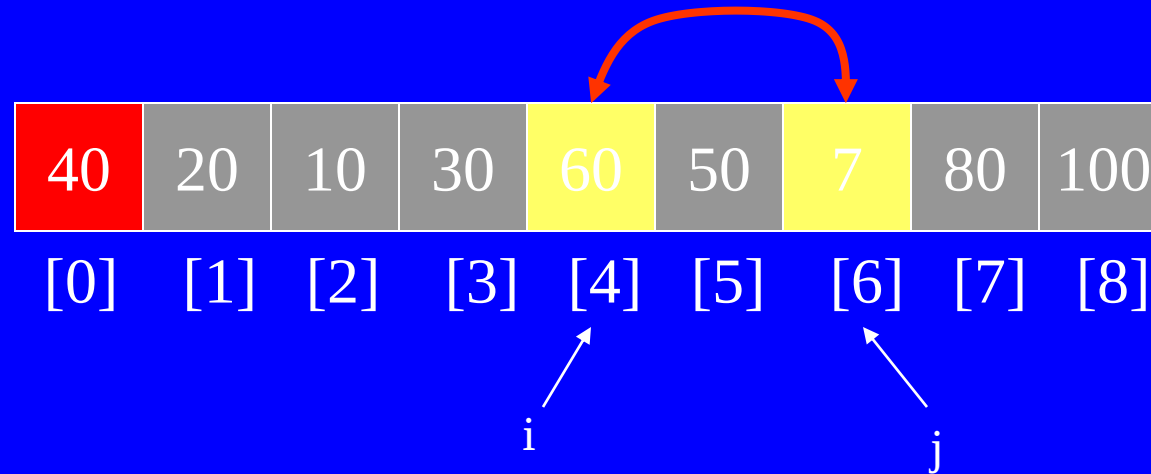
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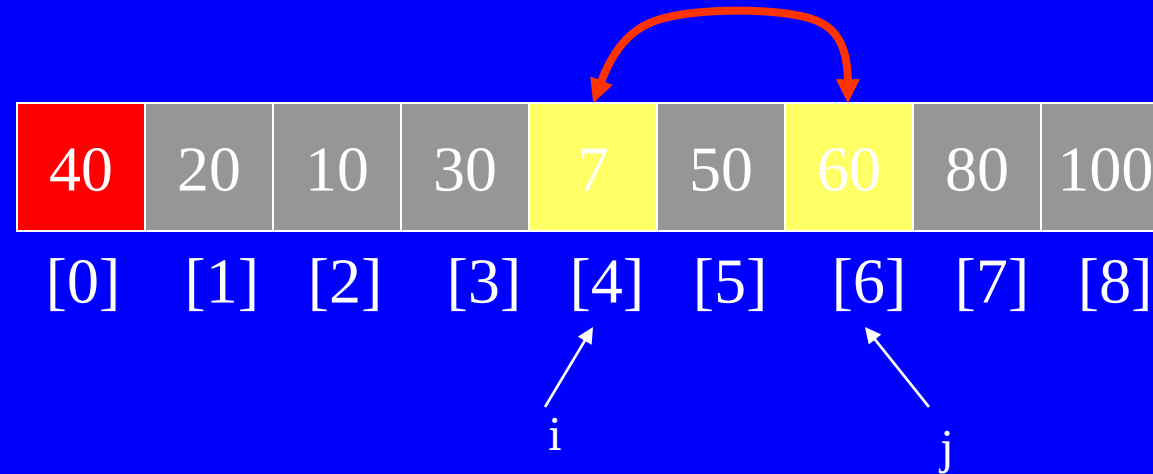
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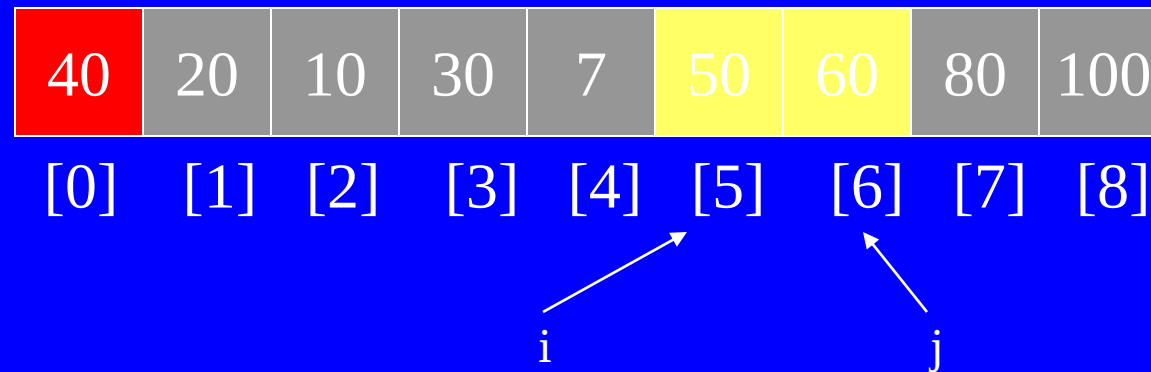
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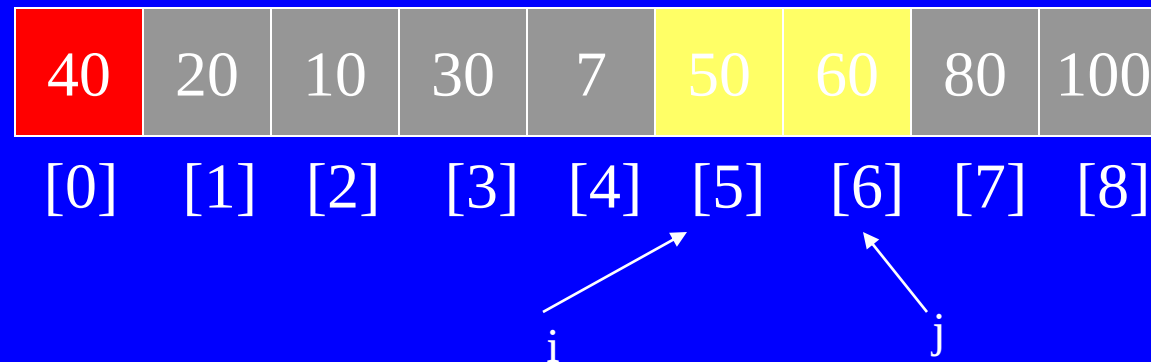
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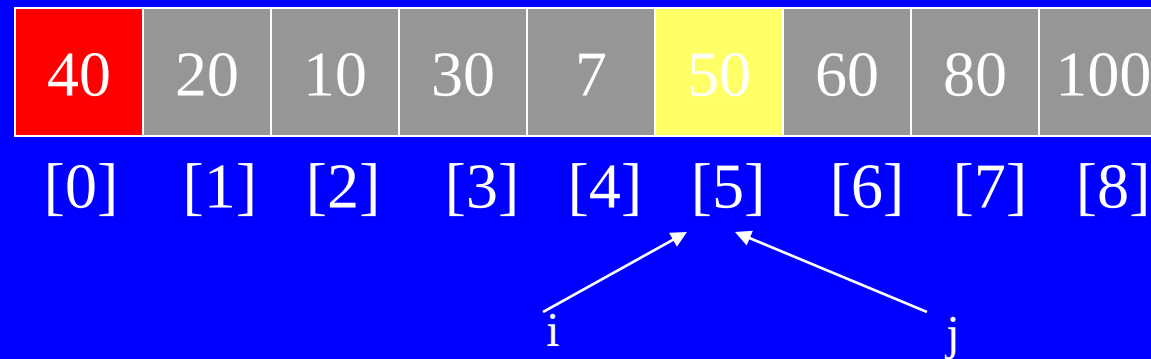
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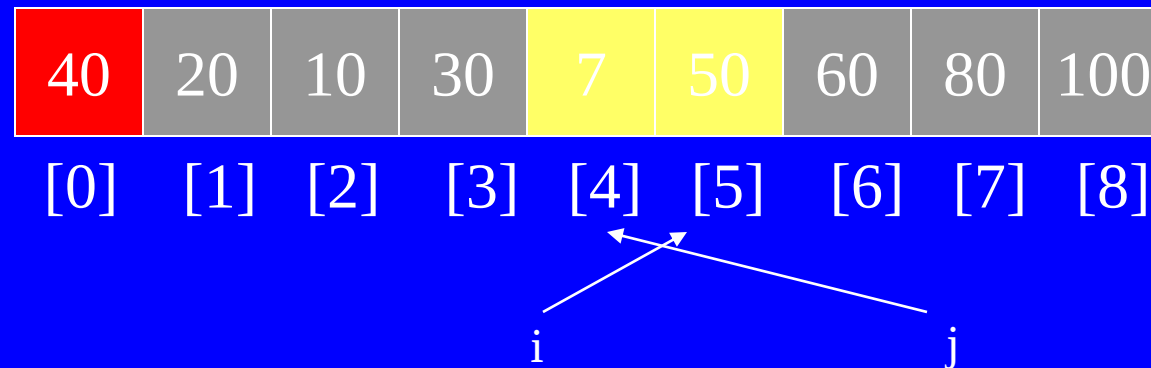
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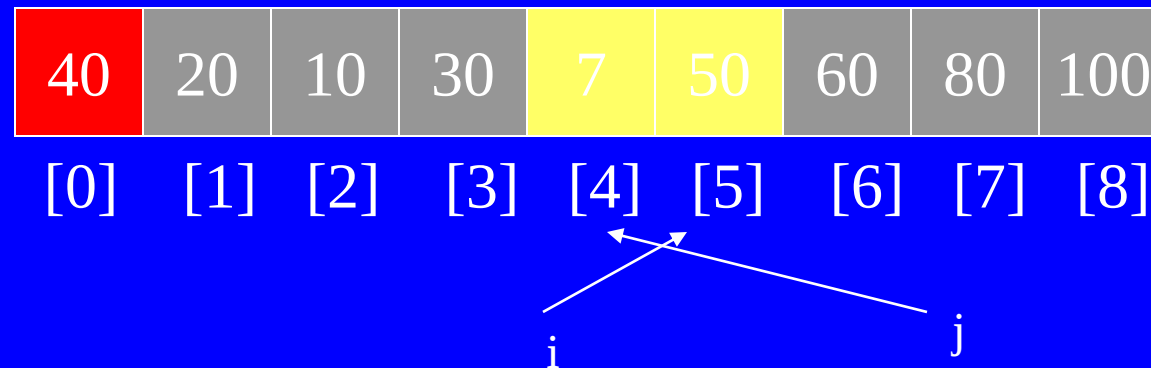
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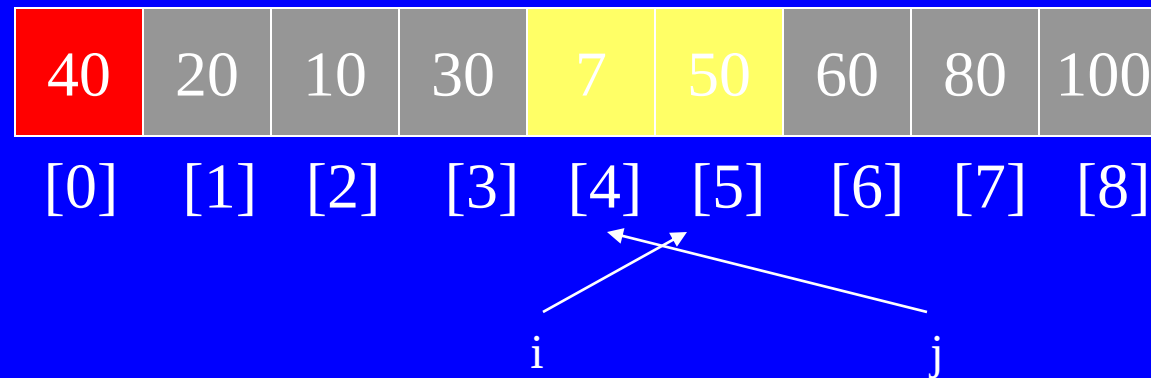
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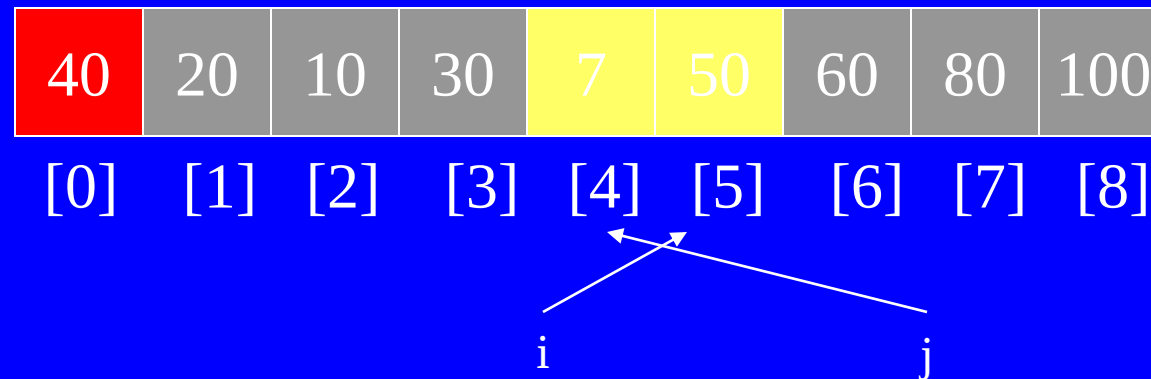
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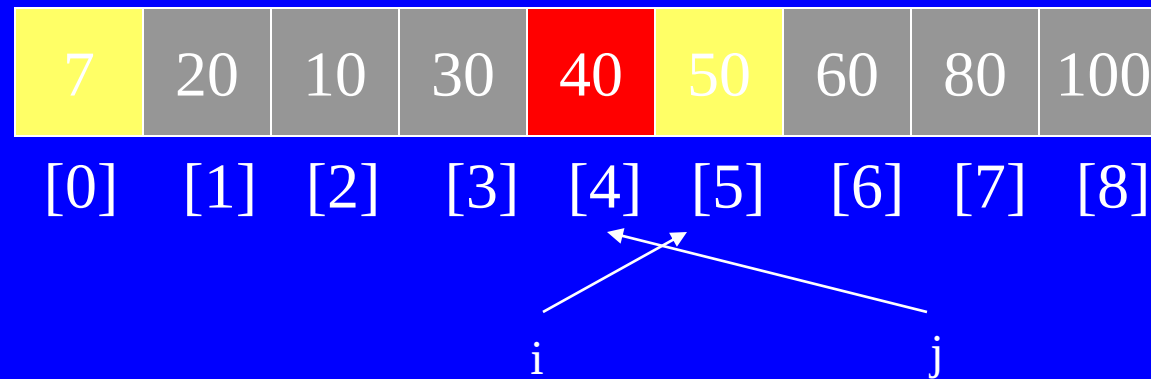
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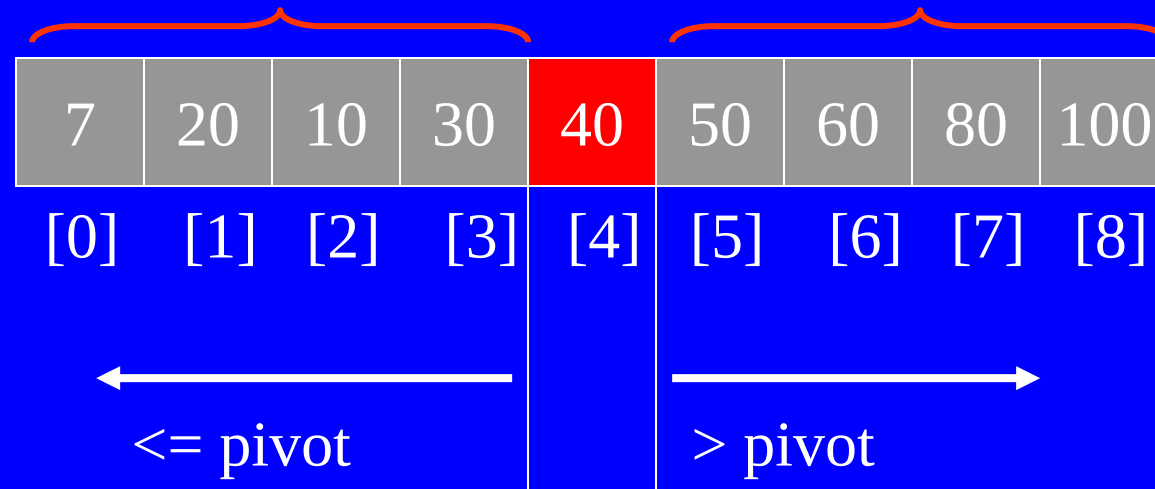
pivot = 40



Partition Result

7	20	10	30	40	50	60	80	100
[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]
←					→			
≤ pivot					> pivot			

Recursion: Quicksort Sub-arrays



Quicksort Analysis

- Worst case: pivot is either maximum or minimum
- One partition is empty, other has size $n-1$
- $T(n) = T(n-1) + n = O(n^2)$
- This will happen for already sorted array
- But mathematically, worst case is very rare
- So most of the time, $T(n) = 2T(n/2) + n = O(n \log n)$
- If at all happens, its because of the fixed choice of the pivot
- Randomized quicksort
- Ideal value of pivot $\rightarrow$ median, divide array into equal halves
- In practice, quicksort is very fast
- In spreadsheet/Python
- No additional list needed, in-place, unlike merge sort

Thank you!!