



**School of Computer Science and Engineering**

**Winter Semester 2023-24**

**Continuous Assessment Test – II**

**SLOT: B2+TB2**

**Programme Name & Branch: B.Tech**

**Course Name & Code: Theory of Computation & BCSE304L**

**Class Number (s): VL2023240500789, VL2023240501047, VL2023240501029, VL2023240500774, VL2023240500845, VL2023240500760, VL2023240501015, VL2023240500784, VL2023240501045, VL2023240500861, VL2023240501036, VL2023240500763, VL2023240501026, VL2023240500766, VL2023240500771, VL2023240501043, VL2023240500797, VL2023240501049, VL2023240501037, VL2023240501042, VL2023240501051, VL2023240501014.**

**Faculty Name (s): NAVAMANI T M, KONATHAM, SUMALATHA, KANAGARAJ R, KANNADASAN R, MADIAJAGAN M, SATHIYA KUMAR C, MOHANA CM, KRISHNA RANI, SAMAL K, UMAMAHESWARI M, SARITHA MURALI, ANAND BIHARI, ANAND M, KRISHNARAJ N, ARUMUGA, ARUN R, SHALINI L, SARWESH P, RAJARAJAN G, SABYASACHI KAMILA, SOMASUNDARAM S K, HUSSAIN AHMED CHOWDHURY, UMA PRIYA D, VISHNUPRIYA.**

**Exam Duration: 90 Min.**

**Maximum Marks: 50**

Q. No.	Question	Max Marks	CO	BL
1.	a) The derivative of a language for a string $x$ is defined as follows: $L_x = \{ y \mid xy \in L \}$ . 1) what is $L_{aa}$ for the language $L = \{ a^{n+1}b^n \mid n \geq 1 \}$ . (1mark) 2) what is $L_b$ for the language $L = \{ a^{n+1}b^n \mid n \geq 1 \}$ . (1 mark) 3) Prove that when $L$ is regular, the derivative of $L$ (ie., $L_x$ ) is also regular. (3 marks)	5	CO3	BL3
	b) Consider the language $L_1 = \{ a^p b^q a^r \mid pq, r \geq 1 \text{ and } p+q \neq r \}$ . Prove or disprove that the language $L$ is regular.	5	CO3	BL2
2.	a) Design a CFG for the language $\gamma = \{ a^n b^m a^{2n} \mid n, m \geq 0 \}$ .	5	CO3	BL2
	b) Consider the grammar $S \rightarrow 0S1 \mid 1S0 \mid \lambda$ , where $S$ is the start symbol, $\{0,1\}$ are terminals. You can see that for every 0 there is a 1 and vice versa in the rule (other than $S \rightarrow \lambda$ ). Is the grammar generates the language $L_2 = \{ w \mid  w _0 =  w _1 \}$ where $ w _a$ refers to the occurrences of $a$ in $w$ . If not, justify with a string in the language $L_2$ that cannot be produced by the given grammar rules. What further changes you need to do in the given grammar in order to generate $L_2$ .	5	CO3	BL2
3.	a) Construct the CYK algorithm for the following CFG and check whether the string “ <b>aabc</b> ” is derivable from the grammar. The Start Symbol is $A$ .	10	CO3	BL3

	<b><math>A \rightarrow CB</math></b> <b><math>B \rightarrow DC \mid b</math></b> <b><math>C \rightarrow BD \mid a</math></b> <b><math>D \rightarrow AD \mid c</math></b>			
4.	a) Construct the given Context Free Grammar <b><math>G=(V,T,P,S)</math></b> into Chomsky normal form. P is the start symbol. <b><math>P \rightarrow RPT</math></b> <b><math>Q \rightarrow RP/Ra</math></b> <b><math>R \rightarrow bRPR \mid b \mid \epsilon</math></b> <b><math>T \rightarrow PaP \mid R \mid aa</math></b>	5	CO3	BL3
	b) Derive the CFG for the given Context Free Language <b><math>L = \{v^r i^{r+s} t^s \mid r, s \geq 0\}</math></b> and check the ambiguity of the given grammar.	5	CO3	BL3
5.	The language L have a strings of balanced “#” symbol and “@” symbol : every “#” symbol can be paired with a unique subsequent of “@” symbol , and every “@” symbol can be paired with a unique preceding “#” symbol. Moreover, the string between any such pair has the same property. Derive the language for the given condition and Construct the PDA for the derived language.	10	CO4	BL4