

**VIT**Vellore Institute of Technology
(Deemed to be University under section 3 of UGC Act, 1956)

REG.NO.:

SCHOOL OF COMPUTER SCIENCE AND ENGINEERING

SLOT: E1 + TE1

CONTINUOUS ASSESSMENT TEST - I
WINTER SEMESTER 2024-2025

Programme Name & Branch : B.Tech & Computer Science and Engineering
Course Code and Course Name : BCSE309L & Cryptography and Network Security
Faculty Name(s) : Applicable to all
Class Number(s) : Applicable to all
Date of Examination : 31.01.25
Exam Duration : 90 minutes **Maximum Marks: 50**

General instruction(s):

- Answer All Questions
- M - Max mark; CO - Course Outcome; BL - Blooms Taxonomy Level (1 - Remember, 2 - Understand, 3 - Apply, 4 - Analyse, 5 - Evaluate, 6 - Create)
- CO1: To know the fundamental mathematical concepts related to security.
- CO2: To understand concept of various cryptographic techniques.

Q. No	Question	M	CO	BL
1.	a) Find $15^{276} \bmod 97$ using the method of fast modular exponentiation. Also, determine the result step-by-step and explain how the exponentiation and modular operations are applied iteratively.	7	CO1	BL3
	b) Compute $\phi(8820)$ using the property of Euler's Totient Function. Describe your steps clearly.	3	CO1	BL3
2.	a) Solve the following congruent equations and find the X value using Chinese Remainder Theorem. $X \equiv 5 \pmod{18}$ $X \equiv 3 \pmod{11}$ $X \equiv 0 \pmod{5}$	7	CO2	BL3
	b) Write the algorithm steps of Miller-Rabin algorithm.	3	CO2	BL3
3.	Explain in detail the state and key initialization, initial state permutation, key stream generation and encryption process of RC4 algorithm with diagram and algorithm steps.	10	CO2	BL2
4.	Construct the shift row matrix from the given sub byte matrix (given in Hex) in AES algorithm. Also, find X and Y values in the given output of mix column matrix using the given fixed matrix and determined shift row matrix (Hint: when finding X & Y values, consider fixed matrix as first matrix and shift row matrix as second matrix)	10	CO2	BL3
	$\begin{bmatrix} 02 & 03 & 01 & 01 \\ 01 & 02 & 03 & 01 \\ 01 & 01 & 02 & 03 \\ 03 & 01 & 01 & 02 \end{bmatrix}$ Fixed matrix $\begin{bmatrix} 63 & 09 & cd & ba \\ 53 & d0 & 30 & 20 \\ 72 & ca & 94 & 6d \\ 04 & 84 & 85 & 25 \end{bmatrix}$ Sub byte matrix $\begin{bmatrix} X & - & - & - \\ Y & - & - & - \\ - & - & - & - \\ - & - & - & - \end{bmatrix}$ Mix column matrix			
5.	Describe the operation of S-Box group and S-Box rule. Also, Find the output of S box group from the output of the whitener (XOR) operation of DES function using given S-Box tables. The output of the whitener is 48-bit values, which serves as the input to the S-boxes. The 48-bit values are given as follows: 110011101010011001110110000011111000101011011110	10	CO2	BL4

Table 6.3 *S-box 1*

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	14	04	13	01	02	15	11	08	03	10	06	12	05	09	00	07
1	00	15	07	04	14	02	13	10	03	06	12	11	09	05	03	08
2	04	01	14	08	13	06	02	11	15	12	09	07	03	10	05	00
3	15	12	08	02	04	09	01	07	05	11	03	14	10	00	06	13

Table 6.4 *S-box 2*

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	15	01	08	14	06	11	03	04	09	07	02	13	12	00	05	10
1	03	13	04	07	15	02	08	14	12	00	01	10	06	09	11	05
2	00	14	07	11	10	04	13	01	05	08	12	06	09	03	02	15
3	13	08	10	01	03	15	04	02	11	06	07	12	00	05	14	09

Table 6.5 *S-box 3*

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	10	00	09	14	06	03	15	05	01	13	12	07	11	04	02	08
1	13	07	00	09	03	04	06	10	02	08	05	14	12	11	15	01
2	13	06	04	09	08	15	03	00	11	01	02	12	05	10	14	07
3	01	10	13	00	06	09	08	07	04	15	14	03	11	05	02	12

Table 6.6 *S-box 4*

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	07	13	14	03	00	6	09	10	1	02	08	05	11	12	04	15
1	13	08	11	05	06	15	00	03	04	07	02	12	01	10	14	09
2	10	06	09	00	12	11	07	13	15	01	03	14	05	02	08	04
3	03	15	00	06	10	01	13	08	09	04	05	11	12	07	02	14

Table 6.7 *S-box 5*

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	02	12	04	01	07	10	11	06	08	05	03	15	13	00	14	09
1	14	11	02	12	04	07	13	01	05	00	15	10	03	09	08	06
2	04	02	01	11	10	13	07	08	15	09	12	05	06	03	00	14
3	11	08	12	07	01	14	02	13	06	15	00	09	10	04	05	03

Table 6.8 *S-box 6*

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	12	01	10	15	09	02	06	08	00	13	03	04	14	07	05	11
1	10	15	04	02	07	12	09	05	06	01	13	14	00	11	03	08
2	09	14	15	05	02	08	12	03	07	00	04	10	01	13	11	06
3	04	03	02	12	09	05	15	10	11	14	01	07	10	00	08	13

Table 6.9 *S-box 7*

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	4	11	2	14	15	00	08	13	03	12	09	07	05	10	06	01
1	13	00	11	07	04	09	01	10	14	03	05	12	02	15	08	06
2	01	04	11	13	12	03	07	14	10	15	06	08	00	05	09	02
3	06	11	13	08	01	04	10	07	09	05	00	15	14	02	03	12

Table 6.10 *S-box 8*

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0	13	02	08	04	06	15	11	01	10	09	03	14	05	00	12	07
1	01	15	13	08	10	03	07	04	12	05	06	11	10	14	09	02
2	07	11	04	01	09	12	14	02	00	06	10	10	15	03	05	08
3	02	01	14	07	04	10	8	13	15	12	09	09	03	05	06	11

① a) $276 = 256 + 16 + 4 = 2^8 + 2^4 + 2^2 \Rightarrow (100010100)$

$15^{276} = 15^{256} \cdot 15^{16} \cdot 15^4 \pmod{97} \Rightarrow 15^{256} \cdot 88 \cdot 89 \pmod{97} = (6600 \pmod{97})$

$15^1 \pmod{97} = 15$

$15^2 \pmod{97} = 225 \pmod{97} = 31$

$15^4 \pmod{97} = 15^2 \cdot 15^2 = 31 \cdot 31 \pmod{97} = 961 \pmod{97} = 88$

$15^8 \pmod{97} = 15^4 \cdot 15^4 = 88 \cdot 88 \pmod{97} = 7744 \pmod{97} = 31$

$15^{16} \pmod{97} = 15^8 \cdot 15^8 = 31 \cdot 31 \pmod{97} = 961 \pmod{97} = 88$

$15^{32} \pmod{97} = 15^{16} \cdot 15^{16} = 88 \cdot 88 \pmod{97} = 7744 \pmod{97} = 31$

$15^{64} \pmod{97} = (15^{32})^2 = 31^2 \pmod{97} = 961 \pmod{97} = 88$

$15^{128} \pmod{97} = 15^{64} \cdot 15^{64} = 88 \cdot 88 \pmod{97} = 7744 \pmod{97} = 31$

$15^{256} \pmod{97} = (15^{128})^2 = 31^2 \pmod{97} = 961 \pmod{97} = 88$

7M

$35 \times 62 \times 88 \pmod{97}$

$2170 \times 88 \pmod{97}$

$36 \cdot 88 \pmod{97}$

$3168 \pmod{97}$

$= 864$

4Q
X = 1C
Y = 5A

⑥ $\phi(8820)$

$$\begin{array}{r} 2 \overline{) 8820} \\ \underline{4410} \\ 2205 \\ \underline{1102} \\ 3735 \\ \underline{5245} \\ 749 \\ \underline{72} \\ 1 \end{array}$$

$$\phi(n) = n \cdot (1 - \frac{1}{p_1}) \cdot (1 - \frac{1}{p_2}) \cdot \dots$$

$$= 8820 \cdot (1 - \frac{1}{2}) \cdot (1 - \frac{1}{3}) \cdot (1 - \frac{1}{5}) \cdot (1 - \frac{1}{7})$$

$$= 8820 \cdot \frac{1}{2} \cdot \frac{2}{3} \cdot \frac{4}{5} \cdot \frac{6}{7} = 2016$$

$M = m_1 \times m_2 \times m_3$
 $18 \cdot 11 \cdot 5 = 990$

$M_1 = \frac{M}{m_1} = \frac{990}{18} = 55$

$M_2 = \frac{M}{m_2} = \frac{990}{11} = 90$

$M_3 = \frac{M}{m_3} = \frac{990}{5} = 198$

② a) $X \equiv 5 \pmod{18}$
 $X \equiv 3 \pmod{11}$
 $X \equiv 0 \pmod{5}$

$a_1 = 5 \quad m_1 = 18$
 $a_2 = 3 \quad m_2 = 11$
 $a_3 = 0 \quad m_3 = 5$

relatively prime

$\gcd(18, 11) = 1$
 $\gcd(18, 5) = 1$
 $\gcd(11, 5) = 1$

$$\begin{array}{r} 11 \overline{) 18} \\ \underline{11} \\ 7 \\ \underline{7} \\ 0 \end{array}$$

$$\begin{array}{r} 11 \overline{) 18} \\ \underline{11} \\ 7 \\ \underline{7} \\ 0 \end{array}$$

$$\begin{array}{r} 5 \overline{) 18} \\ \underline{15} \\ 3 \\ \underline{3} \\ 0 \end{array}$$

$$\begin{array}{r} 5 \overline{) 11} \\ \underline{5} \\ 6 \\ \underline{6} \\ 0 \end{array}$$

$m_1^{-1} \pmod{m_1}$
 $M_1^{-1} = 55^{-1} \pmod{18}$
 $= 55 \cdot x \pmod{18}$

$m_2^{-1} \pmod{m_2}$
 $M_2^{-1} = 90^{-1} \pmod{11}$
 $= 90 \cdot x \pmod{11}$

$m_3^{-1} \pmod{m_3}$
 $M_3^{-1} = 198^{-1} \pmod{5}$
 $= 198 \cdot x \pmod{5}$

$$X = a_1 M_1 M_1^{-1} + a_2 M_2 M_2^{-1} + a_3 M_3 M_3^{-1} \pmod{M}$$

$$= 5 \cdot 55 \cdot 1 + 3 \cdot 90 \cdot 6 \pmod{990}$$

$$= (275 + 1620) \pmod{990}$$

$$= 1895 \pmod{990}$$

$$X = 905$$

2⑥ Miller
 Robin
 Alg

3M

③ RC4 Alg

State & key Initialization (Diagram + Algorithm steps + Explanations) = 2.5 M
 Initial state permutation (") = 2.5 M
 Key stream generation (") = 2.5 M
 encryption (") = 2.5 M

④ find Shift row matrix [2M]

$$\begin{bmatrix} 63 & 09 & cd & ba \\ 53 & d0 & 30 & 20 \\ 72 & ca & 94 & 6d \\ 04 & 84 & 85 & 25 \end{bmatrix} \Rightarrow \begin{bmatrix} 63 & 09 & cd & ba \\ d0 & 30 & 20 & 53 \\ 94 & 6d & 72 & ca \\ 25 & 04 & 84 & 85 \end{bmatrix}$$

No change
 1 byte shift CL
 2 "

Sub byte matrix

find X: [4M]

$$\begin{bmatrix} 02 & 03 & 01 & 01 \\ 01 & 02 & 03 & 01 \\ 01 & 01 & 02 & 03 \\ 03 & 01 & 01 & 02 \end{bmatrix} \times \begin{bmatrix} 63 & 09 & cd & ba \\ d0 & 30 & 20 & 53 \\ 94 & 6d & 72 & ca \\ 25 & 04 & 84 & 85 \end{bmatrix} = \begin{bmatrix} X & - & - & - \\ Y & - & - & - \\ - & - & - & - \\ - & - & - & - \end{bmatrix}$$

$$X = 02 \times 63 \oplus 03 \times d0 \oplus 01 \times 94 \oplus 01 \times 25 = 1C$$

$$Y = 5A \quad (01 \times 63 + 02 \times d0 + 03 \times 94 + 01 \times 25) \text{ mod } 256$$

$$02 \times 63 = (0000 \ 0010) \cdot (0110 \ 0011)$$

$$= x \cdot (x^6 + x^5 + x + 1)$$

$$= x^7 + x^6 + x^2 + x$$

$$03 \times d0 = (0000 \ 0011) \cdot (1101 \ 0000)$$

$$= (x+1) \cdot (x^7 + x^6 + x^4)$$

$$= x^8 + x^7 + x^5 + x^4 + x^6 + x^5 + x^4$$

$$= x^8 + x^6 + x^5 + x^4$$

$$01 \times 94 =$$

$$01 \times 25 =$$

⑤ operation = 1M
 S-box rule = 1M
 o/p S box = 8M

① 110011 $\rightarrow$ (11)₂ $\rightarrow$ (2)₁₀ $\rightarrow$ row
 (1001)₂ = (9)₁₀ $\rightarrow$ col
 Result $\rightarrow$ (11)₁₀ $\rightarrow$ (1011)₂

- ② 04 $\rightarrow$ 0100
- ③ 12 $\rightarrow$ 1100
- ④ 14 $\rightarrow$ 1110
- ⑤ 11 $\rightarrow$ 1011

- ⑥ 01 $\rightarrow$ 0001
- ⑦ 04 $\rightarrow$ 0100
- ⑧ 07 $\rightarrow$ 0111

$$1011 \oplus 0100 \oplus 1100 \oplus 1110 \oplus 1011 \oplus 0001 \oplus 0100 \oplus 0111$$

BCSE309L & Cryptography and Network Security – Winter 2024-25
Answer Key (E1 Slot)

Q. No	Question	M	CO	BL
1.	a) Find $15^{276} \bmod 97$ using the method of fast modular exponentiation. Also, determine the result step-by-step and explain how the exponentiation and modular operations are applied iteratively. Answer: $15^{276} \bmod 97$ $276 = 2^0 + 2^1 + 2^2 + 2^3 + 2^4 + 2^5 + 2^6 + 2^7 + 2^8$ $276 = 2^2 + 2^4 + 2^8$ $276 = 4 + 16 + 256$ $15^{276} \bmod 97 = 15^{4+16+256} \bmod 97$ $= (15^4 \bmod 97 \times 15^{16} \bmod 97 \times 15^{256} \bmod 97) \bmod 97$ $15^2 \bmod 97 = 225 \bmod 97 = 31$ $15^4 \bmod 97 = 31^2 \bmod 97 = 961 \bmod 97 = 88$ $15^8 \bmod 97 = 88^2 \bmod 97 = 7744 \bmod 97 = 81$ $15^{16} \bmod 97 = 81^2 \bmod 97 = 62$ $15^{32} \bmod 97 = 62^2 \bmod 97 = 61$ $15^{64} \bmod 97 = 61^2 \bmod 97 = 35$ $15^{128} \bmod 97 = 35^2 \bmod 97 = 61$ $15^{256} \bmod 97 = 61^2 \bmod 97 = 35$ $15^{276} \bmod 97 = [88 \times 62 \times 35] \bmod 97 = 190960 \bmod 97 = 64$	7		

b) Compute $\phi(8820)$ using the property of Euler's Totient Function. Describe your steps clearly.

Answer:

Property formulae: 1 mark

Detailed steps and correct answer: 2 marks

$\phi(8820)$
 divide by least prime factor

2		8820	
2		4410	
3		2205	
3		735	
5		245	
7		49	
7		7	
		1	

$8820 = 2^2 \times 3^2 \times 5^1 \times 7^2$

Apply the following property

$$\phi(n) = n \times \left(1 - \frac{1}{p_1}\right) \times \left(1 - \frac{1}{p_2}\right) \times \dots$$

$$\phi(8820) = 8820 \times \left[1 - \frac{1}{2}\right] \times \left[1 - \frac{1}{3}\right] \times \left[1 - \frac{1}{5}\right] \times \left[1 - \frac{1}{7}\right]$$

$$= 8820 \times \frac{1}{2} \times \frac{2}{3} \times \frac{4}{5} \times \frac{6}{7}$$

$\phi(8820) = 2016$

3

2. a) Solve the following congruent equations and find the X value using Chinese Remainder Theorem.

$$X \equiv 5 \pmod{18}$$

$$X \equiv 3 \pmod{11}$$

$$X \equiv 0 \pmod{5}$$

Answer:

Relatively prime check, M, M1, M2 & M3 - 3 marks

M1, M2, M3 Inverse values, X value - 4 marks

7

Formed Equations are,

$$1) X \equiv 5 \pmod{18}$$

$$2) X \equiv 3 \pmod{11}$$

$$3) X \equiv 0 \pmod{5}$$

$$a_1 = 5, a_2 = 3, a_3 = 0$$

$$m_1 = 18, m_2 = 11, m_3 = 5$$

Check Relatively Prime/not

$$\text{GCD}[(18, 11), (18, 5), (11, 5)] = 1$$

$$1) 18 = 1 \times 11 + 7$$

$$11 = 1 \times 7 + 4$$

$$7 = 1 \times 4 + 3$$

$$4 = 1 \times 3 + 1 \checkmark$$

$$3 = 3 \times 1 + 0$$

$$2) 18 = 3 \times 5 + 3$$

$$5 = 1 \times 3 + 2$$

$$3 = 1 \times 2 + 1 \checkmark$$

$$2 = 2 \times 1 + 0$$

$$3) 11 = 2 \times 5 + 1 \checkmark$$

$$5 = 5 \times 1 + 0$$

All are relatively prime

Step 1:-

$$M = m_1 \times m_2 \times m_3$$

$$M = 18 \times 11 \times 5$$

$$M = 990$$

Step 2:-

$$M_1 = \frac{M}{m_1} = \frac{990}{18}$$

$$M_1 = 55$$

$$M_2 = \frac{M}{m_2} = \frac{990}{11}$$

$$M_2 = 90$$

Since a_3 is 0,
No need to calculate
 M_3

Step 3:-

$$M_1^{-1} = M_1^{-1} \pmod{m_1}$$

$$= 55^{-1} \pmod{18}$$

$$= 55 \times ? \pmod{18} = 1$$

$$= 55 \times 1 \pmod{18} = 1$$

$$= 55 \pmod{18} = 1$$

$$M_1^{-1} = 1$$

$$M_2^{-1} = M_2^{-1} \pmod{m_2}$$

$$= 90^{-1} \pmod{11}$$

$$= 90 \times ? \pmod{11} = 1$$

$$= 90 \times 6 \pmod{11} = 1$$

$$= 540 \pmod{11} = 1$$

$$M_2^{-1} = 6$$

M_3^{-1} - No need to
Calculate.

	<u>2</u> <u>Step 4:-</u> $X = (a_1 M_1 M_1^{-1} + a_2 M_2 M_2^{-1}) \bmod M$ $= (5 \times 55 \times 1 + 3 \times 90 \times 6) \bmod 990$ $= (275 + 1620) \bmod 990$ $= 1895 \bmod 990$ $X = 905$			
	b) Write the algorithm steps of Miller-Rabin algorithm. Answer: The below answer is one of the algorithm's steps. Some other alternate algorithm steps also there. You can follow any algorithm. Miller-Rabin Primality Test Algorithm Step 1: Find $n-1 = 2^k \times m$ Step 2: Choose 'a' such that $1 < a < n-1$ Step 3: Compute $b_0 = a^m \pmod n, \dots, b_i = b_{i-1}^2 \pmod n$ +1 → Composite -1 → Probably Prime	3		
3.	Explain in detail the state and key initialization, initial state permutation, key stream generation and encryption process of RC4 algorithm with diagram and algorithm steps. Answer: state and key initialization (Diagram + Algorithm steps + Explanation) – 2.5 marks initial state permutation (Diagram + Algorithm steps + Explanation) – 2.5 marks key stream generation (Diagram + Algorithm steps + Explanation) – 2.5 marks encryption process (Diagram + Algorithm steps + Explanation) – 2.5 marks Diagram: The diagram illustrates the RC4 algorithm process in three main stages: State and key initialization (done only once): Shows a key array $K[0..255]$ and a state array $S[0..255]$. The key is mapped to the state array. Initial state permutation (done only once): Shows the state array being mixed with key bytes and permuted. Key stream generation and encryption: A series of boxes showing the iterative process of permuting state values and generating an 8-bit key stream k. This key stream is then XORed with the plaintext P (8 bits) to produce the ciphertext C (8 bits). This process is shown for the first byte, second byte, and last byte.	10		

	Algorithm steps: state and key initialization: - for (i = 0 to 255) { - S[i] ← i - K[i] ← Key [i mod KeyLength] } Initial state permutation: <pre> j ← 0 for (i = 0 to 255) { j ← (j + S[i] + K[i]) mod 256 swap (S[i] , S[j]) } </pre> key stream generation: <pre> i ← (i + 1) mod 256 j ← (j + S[i]) mod 256 swap (S [i] , S[j]) k ← S [(S[i] + S[j]) mod 256] </pre> Encryption process: <pre> // Key is ready, encrypt input P C ← P ⊕ k output C </pre>			
4.	Construct the shift row matrix from the given sub byte matrix (given in Hex) in AES algorithm. Also, find X and Y values in the given output of mix column matrix using the given fixed matrix and determined shift row matrix (Hint: when finding X & Y values, consider fixed matrix as first matrix and shift row matrix as second matrix) $\begin{bmatrix} 02 & 03 & 01 & 01 \\ 01 & 02 & 03 & 01 \\ 01 & 01 & 02 & 03 \\ 03 & 01 & 01 & 02 \end{bmatrix}$ Fixed matrix $\begin{bmatrix} 63 & 09 & cd & ba \\ 53 & d0 & 30 & 20 \\ 72 & ca & 94 & 6d \\ 04 & 84 & 85 & 25 \end{bmatrix}$ Sub byte matrix $\begin{bmatrix} X & - & - & - \\ Y & - & - & - \\ - & - & - & - \\ - & - & - & - \end{bmatrix}$ Mix column matrix Answer: Output of Shift Row: 2 marks Finding X value in the output of Mix column: 4 marks Finding Y value in the output of Mix column: 4 marks	10		

AES

1

Sub byte matrix

b3	09	cd	ba
53	d0	30	20
72	ca	94	6d
04	84	85	25

o/p of shift row

No change	b3	09	cd	ba
1 byte shift	d0	30	20	53
2 byte "	94	6d	72	ca
3 byte "	25	04	84	85

Fixed matrix

02	03	01	01
01	02	03	01
01	01	02	03
03	01	01	02

o/p of shift row

b3	09	cd	ba
d0	30	20	53
94	6d	72	ca
25	04	84	85

Step 1 / Row 1

$$02 \times b3 \oplus 03 \times d0 \oplus 01 \times 94 \oplus 01 \times 25$$

$$02 \times b3 = 0000\ 0010 \times 0110\ 0011$$

$$= x [x^6 + x^5 + x + 1]$$

$$= x^7 + x^6 + x^2 + x \quad \text{--- (1)}$$

$$03 \times d0 = 0000\ 0011 \times 1101\ 0000$$

$$= x+1 [x^7 + x^6 + x^4]$$

$$= x^8 + x^7 + x^5 + x^4 + x^6 + x^4$$

$$= x^8 + x^5 + x^5 + x^4$$

$$= x^4 + x^3 + x + 1 + x^5 + x^5 + x^4$$

$$= x^6 + x^5 + x^3 + x + 1 \quad \text{--- (2)}$$

2

$$\begin{aligned} 01 \times 94 &= 0000 \ 0001 \times 1001 \ 0100 \\ &= 1 [x^7 + x^4 + x^2] \\ &= x^7 + x^4 + x^2 \quad \text{--- (3)} \end{aligned}$$

$$\begin{aligned} 01 \times 25 &= 0000 \ 0001 \times 0010 \ 0101 \\ &= 1 [x^5 + x^2 + 1] \\ &= x^5 + x^2 + 1 \quad \text{--- (4)} \end{aligned}$$

$$\begin{aligned} &= \cancel{x^7} + \cancel{x^6} + \cancel{x^5} + \cancel{x^4} \quad \text{--- (1)} \\ &= \cancel{x^6} + \cancel{x^5} + x^3 + \cancel{x^2} + \cancel{x} \quad \text{--- (2)} \\ &= \cancel{x^7} + x^4 + x^2 \quad \text{--- (3)} \\ &= x^5 + x^2 + x \quad \text{--- (4)} \\ &= x^4 + x^3 + x^2 \\ &= 0001 \ 1100 \end{aligned}$$

$$x = 1C$$

Step 2 | Row 2

$$01 \times 63 \oplus 02 \times d0 \oplus 03 \times 94 \oplus 01 \times 25$$

$$01 \times 63 = x^6 + x^5 + x + 1 \quad \text{--- (1)}$$

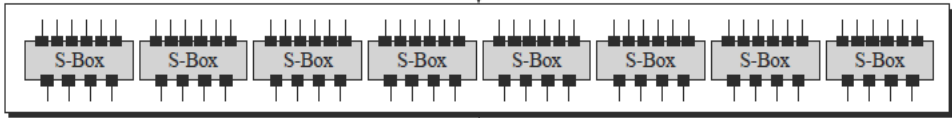
$$02 \times d0 = 0000 \ 0010 \times 1101 \ 0000$$

$$= x [x^7 + x^6 + x^4]$$

$$= x^8 + x^7 + x^5$$

$$= x^4 + x^3 + x + 1 + x^7 + x^5$$

$$= x^7 + x^5 + x^4 + x^3 + x + 1 \quad \text{--- (2)}$$

	$ \begin{aligned} 03 \times 94 &= x+1 \left[x^7+x^4+x^2 \right] \\ &= x^8+x^5+x^3+x^7+x^4+x^2 \\ &= \cancel{x^4}+\cancel{x^3}+x+1+\cancel{x^5}+\cancel{x^3}+\cancel{x^7}+\cancel{x^4}+x^2 \\ &= x^7+x^5+x^2+x+1 \quad \text{--- (3)} \\ 01 \times 25 &= x^5+x^2+1 \quad \text{--- (4)} \\ &= x^6+\cancel{x^5}+\cancel{x}+1 \quad \text{--- (1)} \\ &= \cancel{x^7}+\cancel{x^5}+x^4+x^3+\cancel{x}+1 \quad \text{--- (2)} \\ &= \cancel{x^7}+\cancel{x^5}+\cancel{x^4}+x+x+1 \quad \text{--- (3)} \\ &= x^5+x^2+1 \quad \text{--- (4)} \\ &= x^6+x^4+x^3+x \\ &= 01011010 \\ Y &= 5A \end{aligned} $			
5.	Describe the operation of S-Box group and S-Box rule. Also, Find the output of S box group from the output of the whitener (XOR) operation of DES function using given S-Box tables. The output of the whitener is 48-bit values, which serves as the input to the S-boxes. The 48-bit values are given as follows: 110011101010011001110110000011111000101011011110 Answer: Operation of S-box group: 1 mark S-Box Rule: 1 mark Output of the S box group (Each s-box:1 mark): 8 marks Operation of S-box group: (Either draw this diagram or theoretically, they can explain) 48-bit input  32-bit output Fig. 6.7 S-boxes	10		

S-Box Rule:

(Either draw this diagram or theoretically, they can explain)

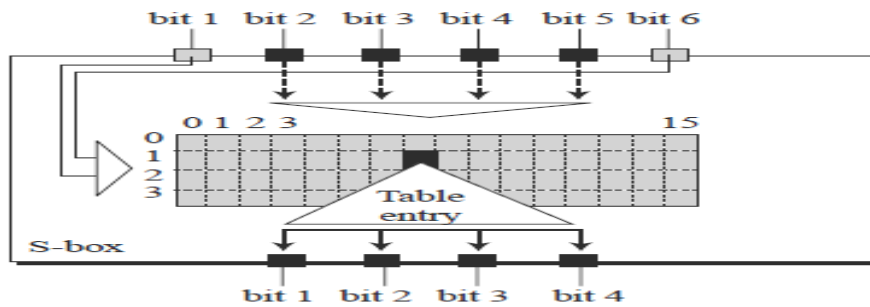


Fig. 6.8 S-box rule

Output of the S-Box Group:

Inputs

<u>110011</u> SBox1	<u>101010</u> S2	<u>011001</u> S3	<u>110110</u> S4
<u>000011</u> S5	<u>111000</u> S6	<u>101011</u> S7	<u>011110</u> S8

SBox1:-

110011 → 11 (Binary) - 3 (Decimal) → Row
 110011 → 1001 (B) - 9 (D) - column
 check in Row 3 & column 9
 Result is 11 (D) → 1011 (B)

SBox2:-

101010 → 10 (B) - 2 (D)
 101010 → 0101 (B) - 5 (D)
 Row 2 & column 5
 Result is 04 (D) → 0100 (B)

SBox3:-

011001 → 01 (B) - 1 (D) → R
 011001 → 1100 (B) - 12 (D) - C
 12 (D) → 1100 (B)

SBox4:-

110110 → 10 (B) - 2 (D) → R
 110110 → 1011 (B) - 11 (D) → C
 14 (D) → 1110 (B)

S Box 5:-

$\underline{000011} \rightarrow 01 (B) - 1 (D) \rightarrow R$
 $0 \underline{00011} \rightarrow 0001 (B) - 1 (D) \rightarrow C$

$11 (D) - \boxed{1011 (B)}$

S Box 6:-

$\underline{111000} \rightarrow 10 (B) - 2 (D) \rightarrow R$
 $\underline{111000} \rightarrow 1100 (B) - 12 (D) \rightarrow C$

$01 (D) - \boxed{0001 (B)}$

S Box 7:-

$\underline{101011} \rightarrow 11 (B) - 3 (D) \rightarrow R$
 $\underline{101011} \rightarrow 0101 (B) - 5 (D) \rightarrow R$

$04 (D) - \boxed{0100 (B)}$

S Box 8:-

$\underline{011110} \rightarrow 00 (B) - 0 (D) \rightarrow R$
 $\underline{011110} \rightarrow 1111 (B) - 15 (D) \rightarrow C$

$07 (D) - \boxed{0111 (B)}$

Final Answer:-

$1011 \ 0100 \ 1100 \ 1110 \ 1011 \ 0001 \ 0100 \ 0111$
