

**VIT**Vellore Institute of Technology
(Deemed to be University under section 3 of UGC Act, 1956)

REG.NO.: 34W100785

SLOT: D2+TD2

SCHOOL OF COMPUTER SCIENCE AND ENGINEERING
CONTINUOUS ASSESSMENT TEST - II
FALL SEMESTER 2025-2026

Programme Name & Branch : M.Tech. (Integrated) – MIC, MID
Course Code and Course Name : CSI1003 – Formal Languages and Automata Theory
Faculty Name(s) : Prof. G Rajarajan, Prof.V Vijayarajan , Prof. R Kannadasan
Class Number(s) : 2465, 2466, 2468
Date of Examination : 08.10.2025 - ForeNoon
Exam Duration : 90 minutes **Maximum Marks: 50**

General instruction(s):

- Answer All Questions
- M - Max mark; CO – Course Outcome; BL – Blooms Taxonomy Level (1 – Remember, 2 – Understand, 3 – Apply, 4 – Analyse, 5 – Evaluate, 6 – Create)
- CO1: Model, compare and analyse different computational models
CO2: Apply rigorously formal mathematical methods to prove properties of languages, grammars and automata.
CO3: Identify limitations of some computational models and possible methods of proving them

Q. No	Question	Marks	CO	BL
1.	a) Let us assume L1 and L2 are Regular Sets. There is an operation called "connect" which connects languages and it is defined as $\{m_a p_k, m_b p_{k-1}, m_c p_{k-2}, \dots, m_{k-1} p_b, m_k p_a \mid m_a, m_b, m_c, \dots, m_{k-1}, m_k \in L1 \text{ and } p_k, p_{k-1}, p_{k-2}, \dots, p_b, p_a \in L2\}$. As L1 and L2 are Regular (as per assumption), will they be closed under the operation "connect"? Justify your answer.	5	2	3
	b) Using pumping lemma of Regular language, either prove or disprove the following: $L = \{ \#^c \mid \text{where } c \text{ is non-composite number over } \Sigma = \{\#\} \}$ is not a Regular Set.	5		
2.	a) Let $L = \{a^n b^n \mid n > 0\}$ is a Context Free Language. Tell whether the grammar generating this language L is inherently ambiguous or unambiguous? If so, brief the difference between the unambiguous language and inherently ambiguous language.	4	3	4
	b) Convert the following into CNF: $S \rightarrow 0X1X$ $X \rightarrow 1X \mid 0Y \mid \epsilon$ $Y \rightarrow 2 \mid X$ NOTE: ϵ means epsilon (empty string) and S is the start Symbol and terminal set = $\{0,1,2\}$. Also, find the number of derivations required to derive a valid string whose length is 1000, from your resultant CNF.	6		



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3.	Use CYK algorithm and check the membership of "baaba" and infer the following table containing cells marked 'X' are True or False for the given grammar. (S is the start symbol of Grammar) $S \rightarrow AB \mid BC$ $A \rightarrow BA \mid a$ $B \rightarrow CC \mid b$ $C \rightarrow AB \mid a$ <table><tr><td></td><td>5</td><td>4</td><td>3</td><td>2</td><td>1</td></tr><tr><td>1</td><td>X</td><td>X</td><td></td><td></td><td>X</td></tr><tr><td>2</td><td></td><td></td><td></td><td></td><td></td></tr><tr><td>3</td><td></td><td></td><td>X</td><td></td><td></td></tr><tr><td>4</td><td>X</td><td></td><td></td><td></td><td></td></tr><tr><td>5</td><td></td><td></td><td></td><td></td><td></td></tr></table> - The cell (1,1) indicates the substring "b" is also member of this Grammar. (True / False) ? - The cell (4,5) indicates the substring "ba" can be derivable from S . (True / False) ? - The cell (1,5) indicates the string "baaba" is derivable from the given grammar. (True / False) ? - The cell (1,4) indicates the substring "baab" cannot be derivable from the start symbol of grammar (True / False)? - The cell (3,3) indicates the substring "a" satisfies the membership of grammar (True / False) ?		5	4	3	2	1	1	X	X			X	2						3			X			4	X					5						10	3	4
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4	X																																							
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4.	Design PDA that recognizes the following language $L = \{ i^p j^q k^r l^s \mid p, q, r, s > 0 \text{ and } p = s \text{ over } \Sigma = \{ i, j, k, l \} \}$. Also verify the strings "iijkl" and "ijjkl" using instantaneous descriptions of PDA	10	1	3																																				
5.	Let $L = \{ \alpha\beta \mid \alpha \in \{0,1\}^+ \text{ and } \beta \text{ is exact copy of } \alpha \}$ Statement 1: Deterministic Pushdown Automata (DPDA) can be designed for the given language and in fact Non-deterministic Pushdown Automata (NPDA) can also be designed for the same. Statement 2: DPDA is not at all possible to design and only Non-Deterministic Pushdown Automata is possible as the power of NPDA is more than DPDA Statement 3: This is a special case in which only DPDA can be designed, but not NPDA. Statement 4: Neither DPDA nor NPDA can be designed for the given language. Out of these 4 statements, which one is correct? Justify Properly with valid explanations or proof, examples to validate your claim.	10	1	5																																				